

## On the exact simulation of functionals of stationary Markov chains

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ID-IMAG Laboratory

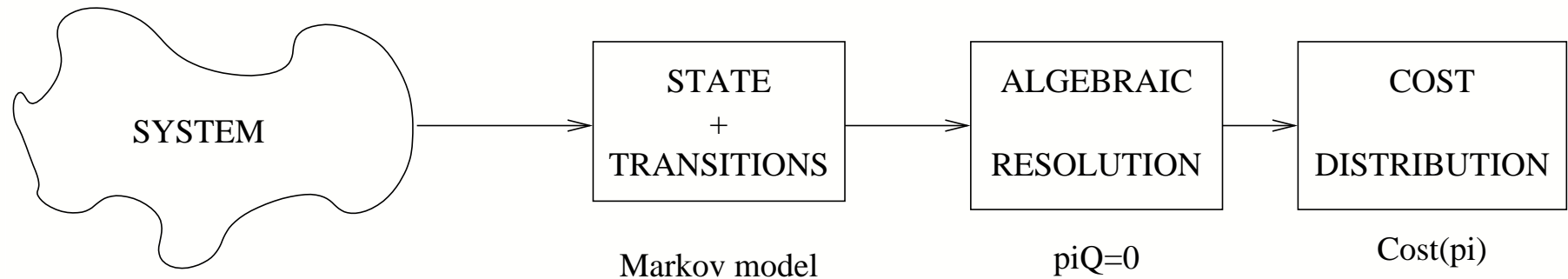


Universities of Grenoble

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# Modeling discrete event systems



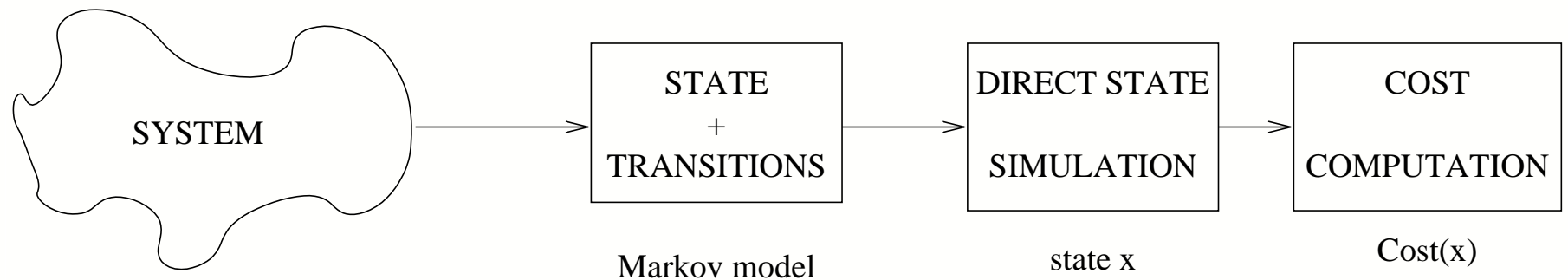
## Difficulties:

- complex structure
- large state space
- analytical/numerical method
- approximation/bounding techniques

⇒ reduction of the state space



# Modeling discrete event systems

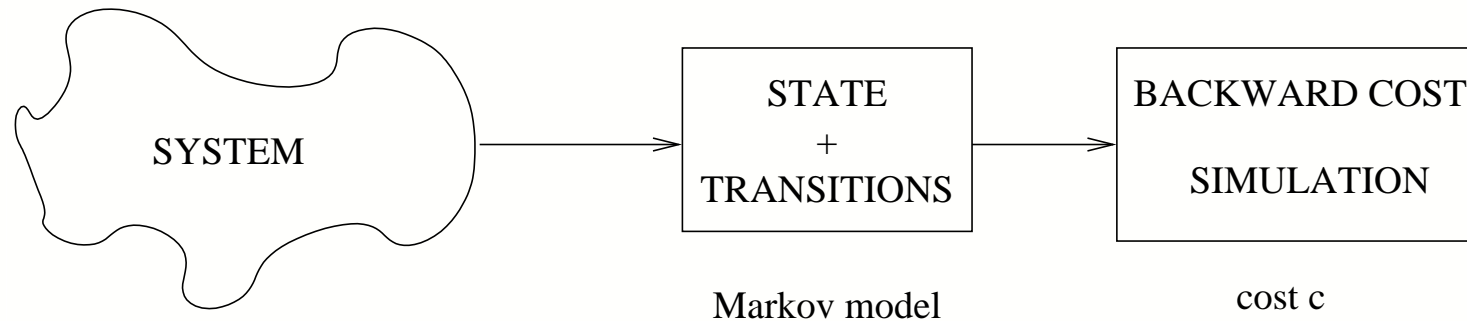


## Difficulties:

- stopping criteria : burn in time
- simulation biases  $\|\pi_n - \pi_\infty\|$
- estimation biases : confidence intervals  $\mathcal{O}\left(\frac{1}{\sqrt{n}}\right)$



# Modeling discrete event systems



Properties:

- Exact stopping criteria  
⇒ no simulation bias

Constraints:

- $N$  parallel trajectories



# Iterated system of functions

$\mathcal{X}$  state space (size  $n$ );  $\mathcal{U}$  set of external input values

Transition function  $\Phi$

$$\begin{aligned}\Phi : \mathcal{X} \times \mathcal{U} &\longrightarrow \mathcal{X} \\ (x, u) &\longmapsto \Phi(x, u)\end{aligned}$$

If  $\{U_n\}_{n \in \mathbb{Z}}$  is IID then

$$X_0 = x_0, \quad X_{n+1} = \Phi(X_n, U_{n+1})$$

is a Markov chain (stochastic recursive sequence).

$$\{\Phi(\cdot, u)\}_{u \in [0, 1[}$$

is an iterated system of function.

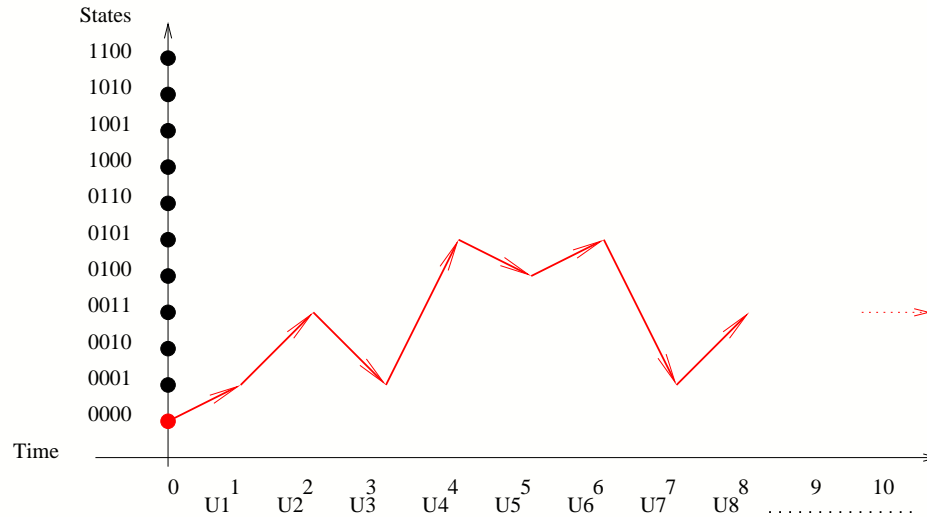
Reciprocally, given a transition matrix  $Q$  it is possible to build a family of function  $\Phi(\cdot, u)$  such that the associated process is a markov chain with transition matrix  $Q$ .

$\implies$  Simulation kernel

## Problem : several representation of $Q$



# Forward simulation



## Forward simulation

$x \leftarrow x_0;$

**repeat**

$u \leftarrow \text{Random};$

$x \leftarrow \Phi(x, u);$

**until** stopping criterium

return  $x$

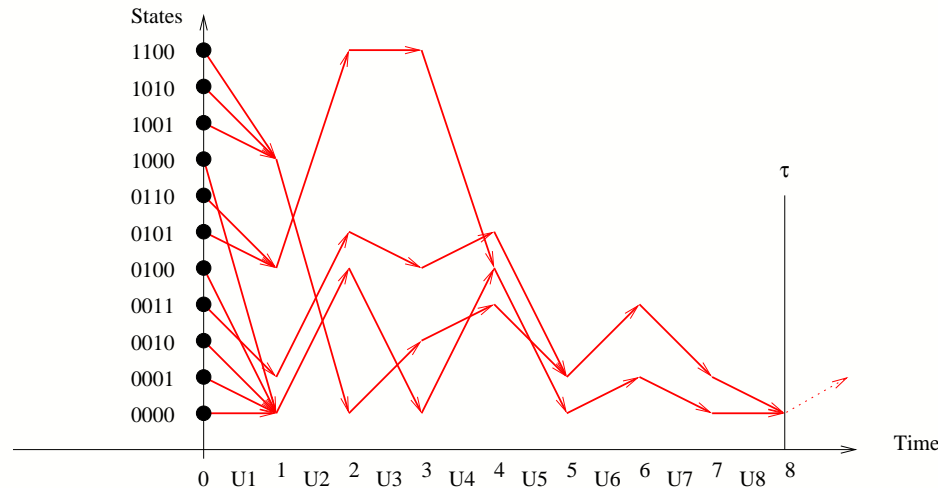
compute statistics

## Problems :

- choice of the initial state
- choice of stopping criterium
- error computation (bias)

# Forward coupling simulation

Forward coupling algorithm



**for all**  $x \in \mathcal{X}$  **do**

$y(x) \leftarrow x$

**end for**

**repeat**

$u \leftarrow \text{Random};$

**for all**  $x \in \mathcal{X}$  **do**

$y(x) \leftarrow \Phi(y(x), u);$

**end for**

**until** All  $y(x)$  are equal

**return**  $y(x)$

**Problems :**

- generated state  $y(x)$

**does not follow the stationary distribution**

Coupling time  $\tau$

$$\tau = \min\{n \in \mathbb{N}, |\Phi(\Phi(\dots(\Phi(\mathcal{X}, U_1), \dots), U_{n-1}), U_n)| = 1\}$$

- Is  $\tau$  almost surely **finite** ?  $\mathbb{P}(\tau < +\infty) = 1$  ?

# Backward coupling simulation

## Idea :

Propp & Wilson(1996)

- reverse time
- run  $N$  parallel trajectories
- wait for coupling.

$$\mathcal{Z}_n = \Phi(\Phi(\cdots(\Phi(\mathcal{X}, U_{-n+1}), \cdots), U_{-1}), U_0).$$

potential set of reachable states at step  $n$

```
for all  $x \in \mathcal{X}$  do
   $y(x) \leftarrow x$ 
end for
repeat
   $u \leftarrow \text{Random};$ 
  for all  $x \in \mathcal{X}$  do
     $y(x) \leftarrow y(\Phi(x, u));$ 
  end for
until All  $y(x)$  are equal
return  $y(x)$ 
```

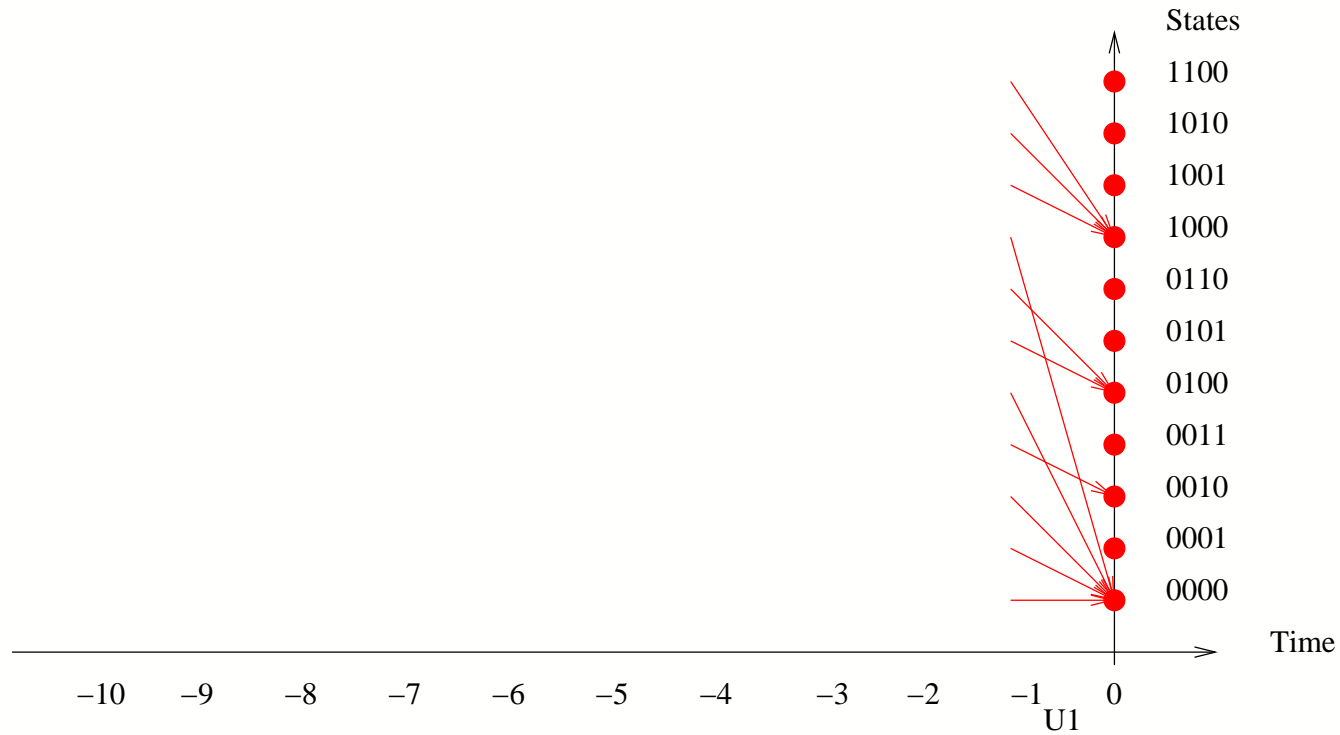




# Backward simulation example



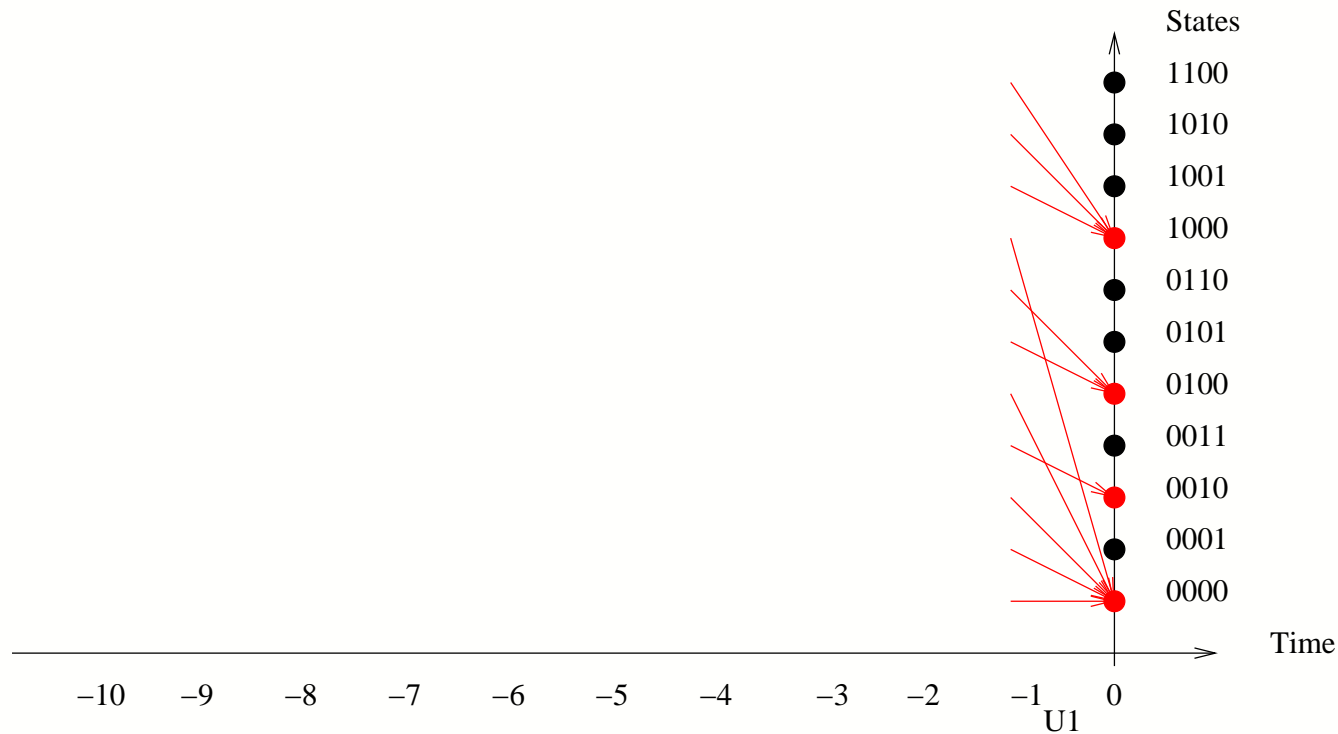
# Backward simulation example



$$Z_0 = \mathcal{X}$$



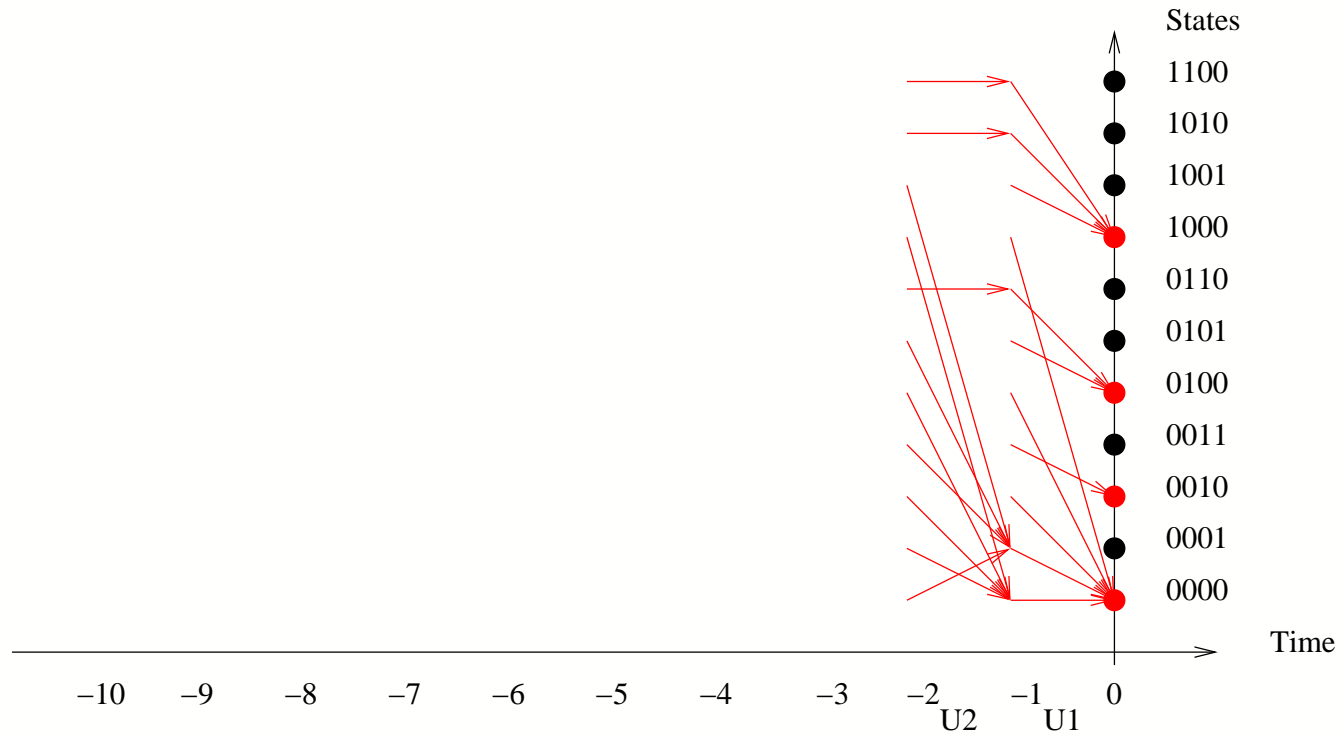
# Backward simulation example



$$\mathcal{Z}_0 = \mathcal{X}$$

$$\mathcal{Z}_1 = \{0000, 0010, 0100, 1000\}$$

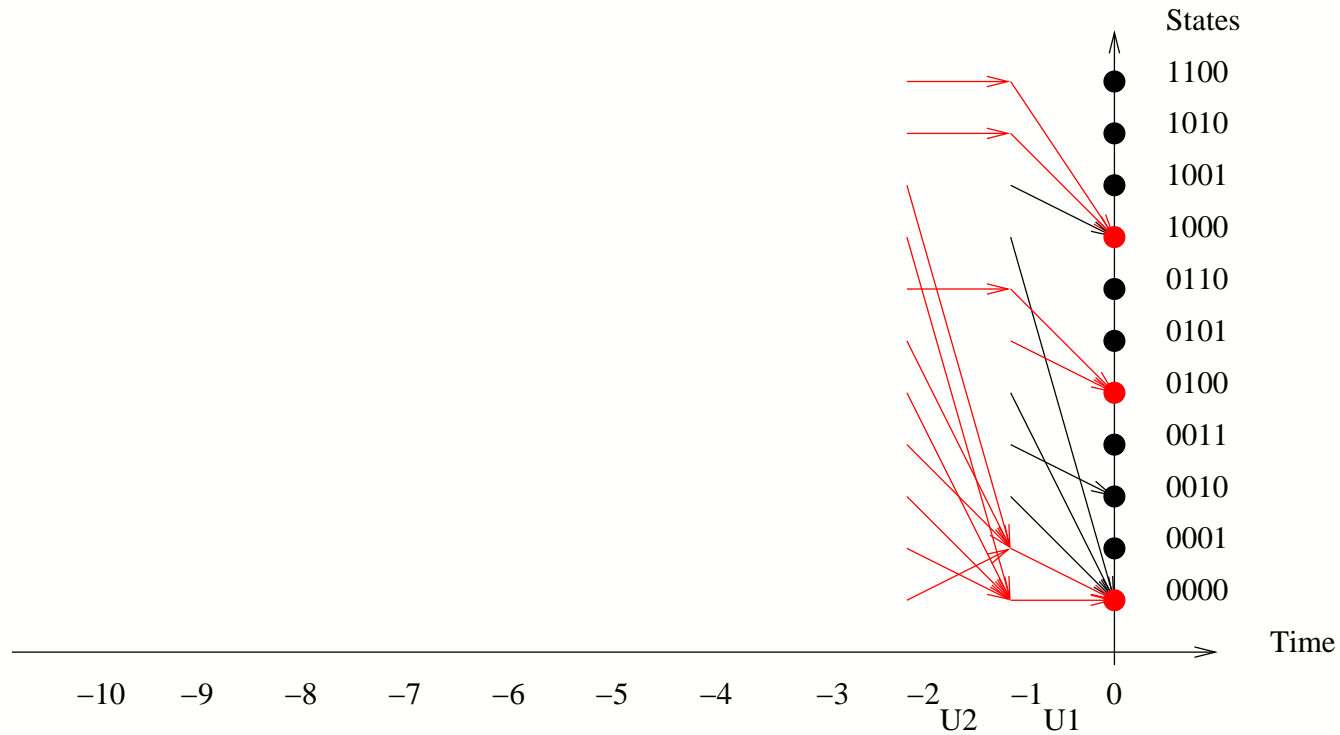
# Backward simulation example



$$\mathcal{Z}_0 = \mathcal{X}$$

$$\mathcal{Z}_1 = \{0000, 0010, 0100, 1000\}$$

# Backward simulation example



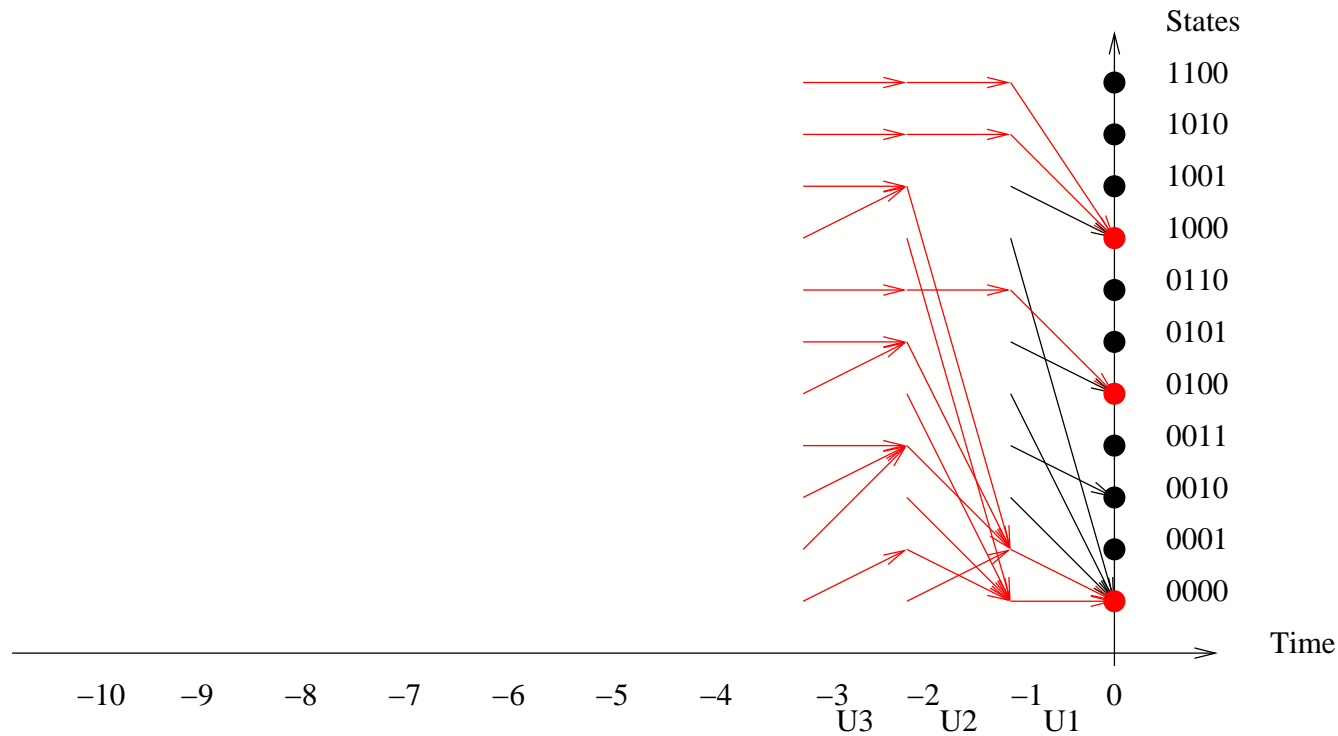
$$\mathcal{Z}_0 = \mathcal{X}$$

$$\mathcal{Z}_1 = \{0000, 0010, 0100, 1000\}$$

$$\mathcal{Z}_2 = \{0000, 0100, 1000\}$$



# Backward simulation example



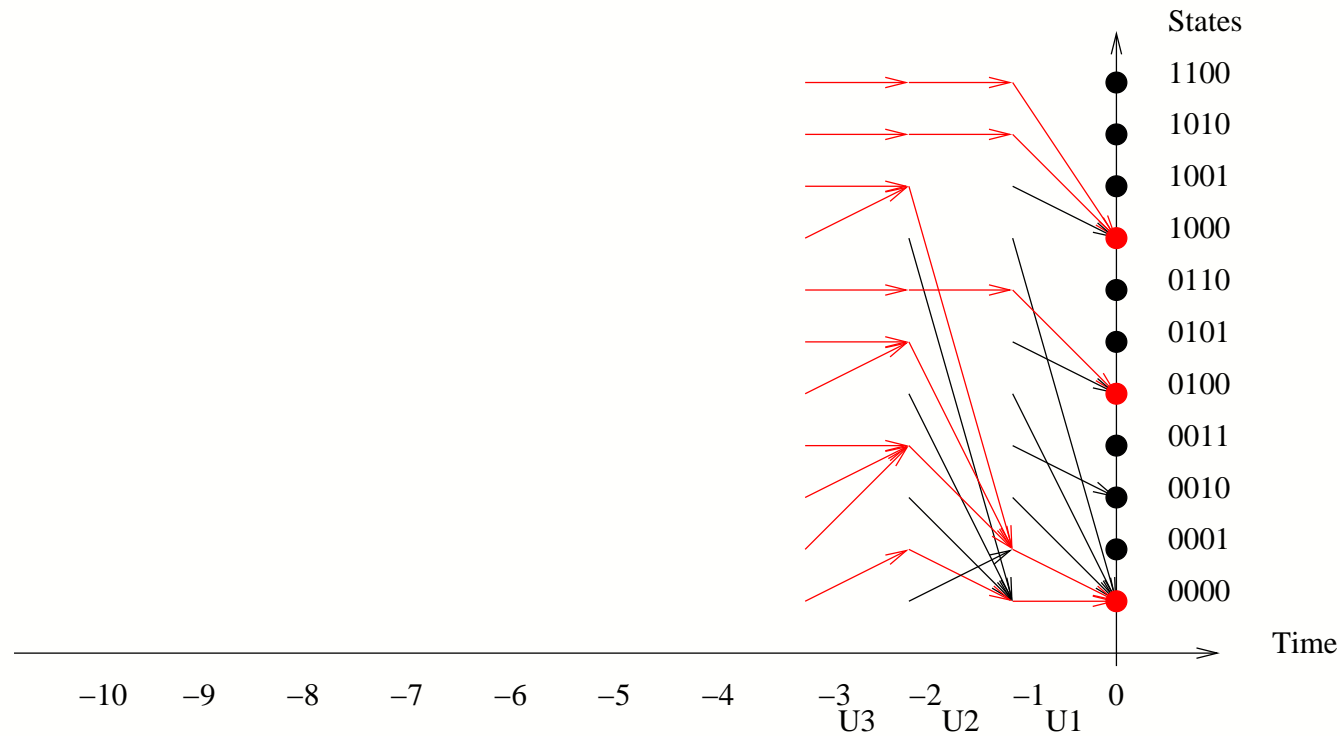
$$\mathcal{Z}_0 = \mathcal{X}$$

$$\mathcal{Z}_1 = \{0000, 0010, 0100, 1000\}$$

$$\mathcal{Z}_2 = \{0000, 0100, 1000\}$$



# Backward simulation example



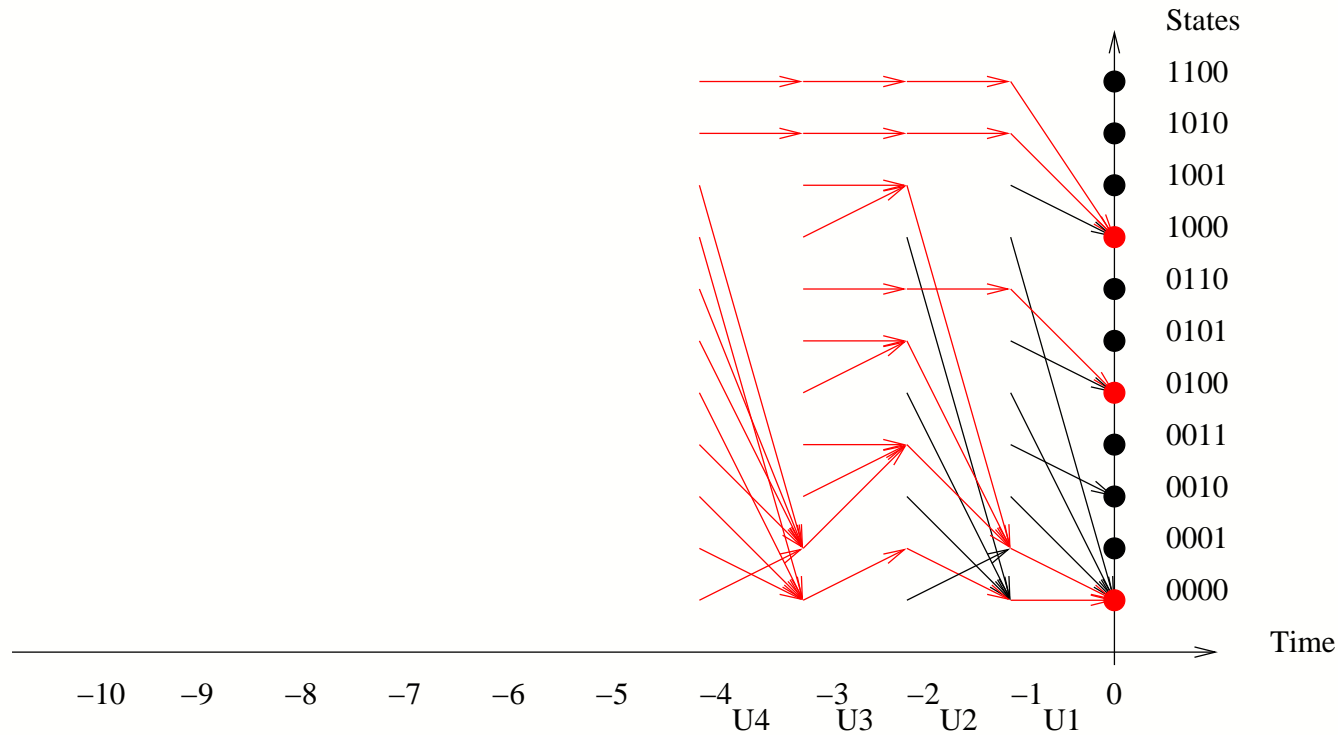
$$\mathcal{Z}_0 = \mathcal{X}$$

$$\mathcal{Z}_3 = \{0000, 0100, 1000\}$$

$$\mathcal{Z}_1 = \{0000, 0010, 0100, 1000\}$$

$$\mathcal{Z}_2 = \{0000, 0100, 1000\}$$

# Backward simulation example



$$\mathcal{Z}_0 = \mathcal{X}$$

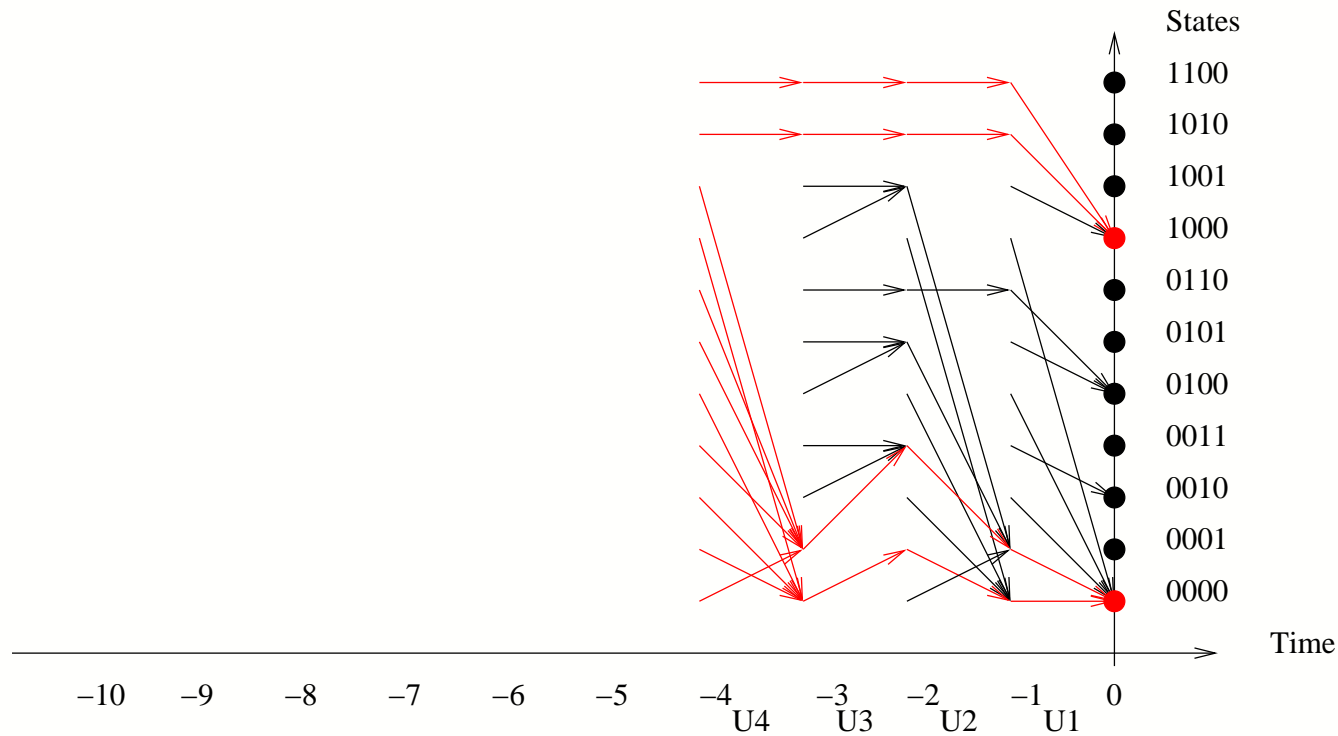
$$\mathcal{Z}_3 = \{0000, 0100, 1000\}$$

$$\mathcal{Z}_1 = \{0000, 0010, 0100, 1000\}$$

$$\mathcal{Z}_2 = \{0000, 0100, 1000\}$$



# Backward simulation example



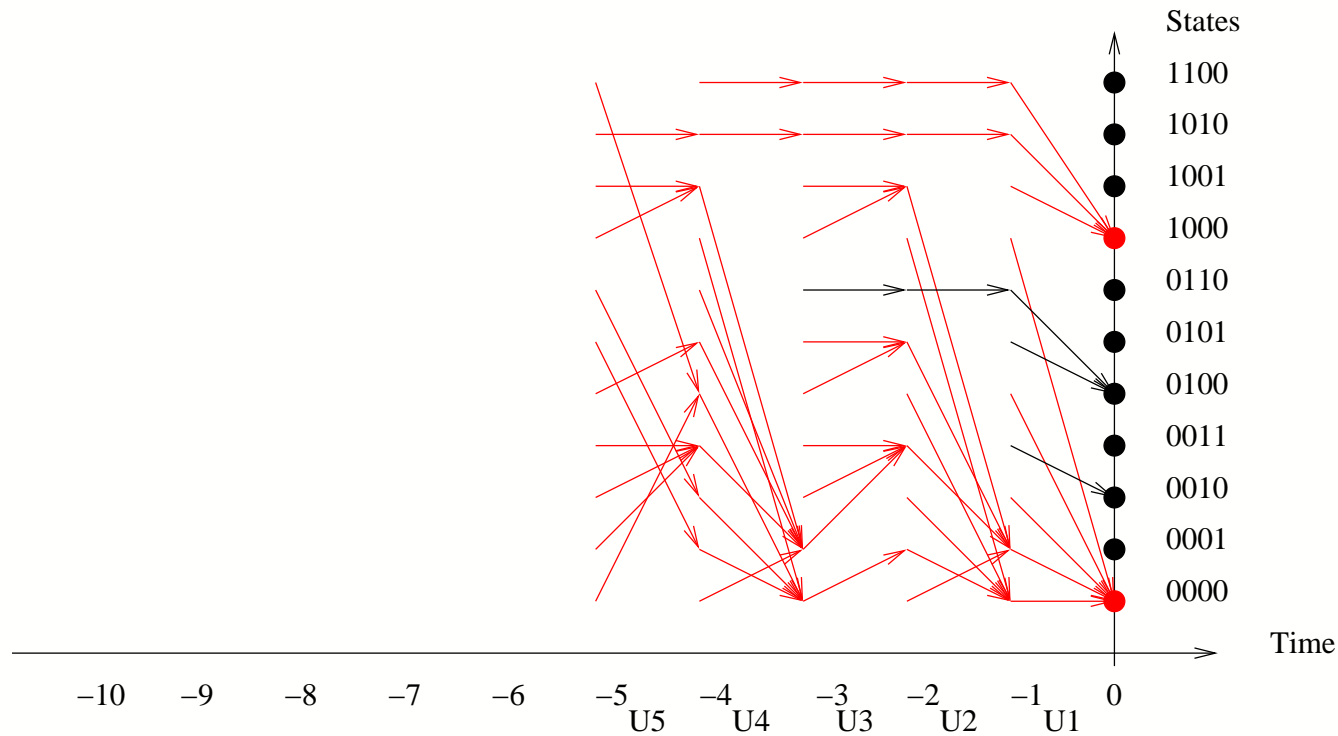
$$Z_0 = \mathcal{X}$$

$$Z_3 = \{0000, 0100, 1000\}$$

$$Z_1 = \{0000, 0010, 0100, 1000\} \quad Z_4 = \{0000, 1000\}$$

$$Z_2 = \{0000, 0100, 1000\}$$

# Backward simulation example



$$\mathcal{Z}_0 = \mathcal{X}$$

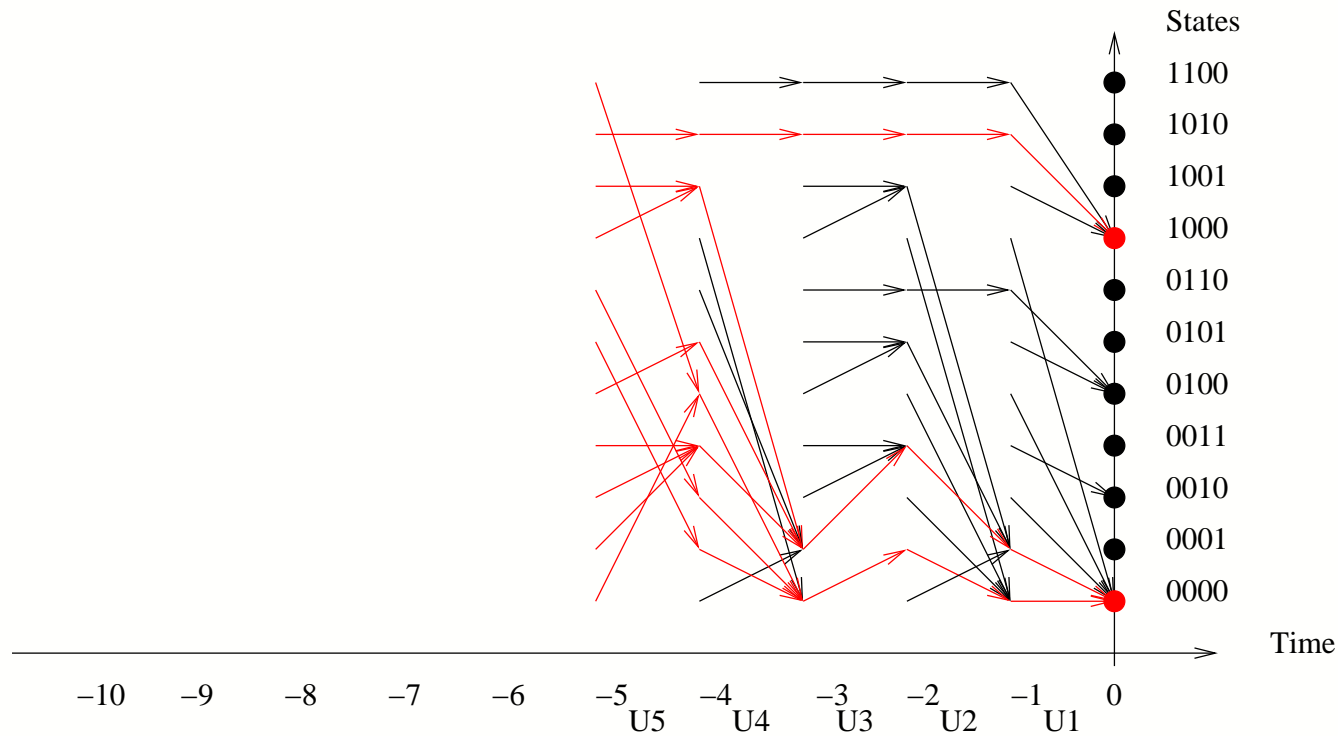
$$\mathcal{Z}_3 = \{0000, 0100, 1000\}$$

$$\mathcal{Z}_1 = \{0000, 0010, 0100, 1000\}$$

$$\mathcal{Z}_4 = \{0000, 1000\}$$

$$\mathcal{Z}_2 = \{0000, 0100, 1000\}$$

# Backward simulation example



$$\mathcal{Z}_0 = \mathcal{X}$$

$$\mathcal{Z}_1 = \{0000, 0010, 0100, 1000\}$$

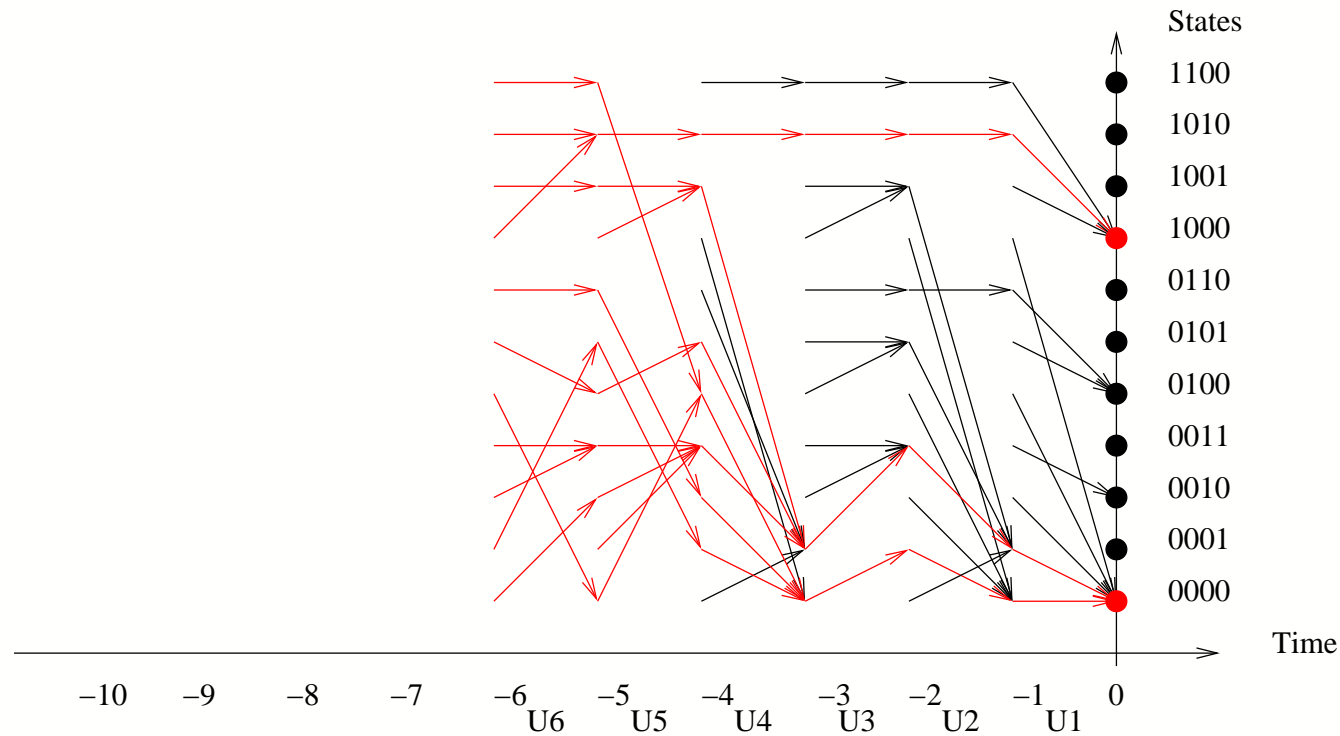
$$\mathcal{Z}_2 = \{0000, 0100, 1000\}$$

$$\mathcal{Z}_3 = \{0000, 0100, 1000\}$$

$$\mathcal{Z}_4 = \{0000, 1000\}$$

$$\mathcal{Z}_5 = \{0000, 1000\}$$

# Backward simulation example



$$\mathcal{Z}_0 = \mathcal{X}$$

$$\mathcal{Z}_1 = \{0000, 0010, 0100, 1000\}$$

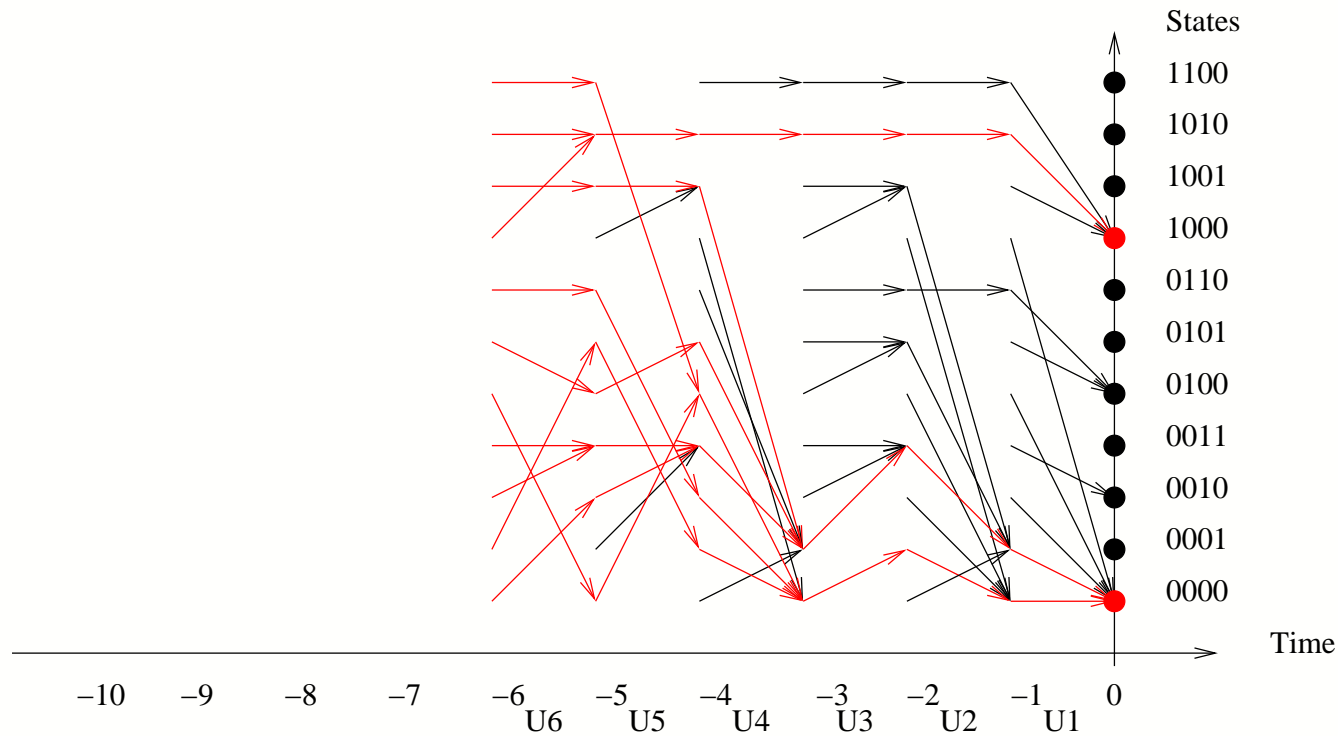
$$\mathcal{Z}_2 = \{0000, 0100, 1000\}$$

$$\mathcal{Z}_3 = \{0000, 0100, 1000\}$$

$$\mathcal{Z}_4 = \{0000, 1000\}$$

$$\mathcal{Z}_5 = \{0000, 1000\}$$

# Backward simulation example



$$\mathcal{Z}_0 = \mathcal{X}$$

$$\mathcal{Z}_1 = \{0000, 0010, 0100, 1000\}$$

$$\mathcal{Z}_2 = \{0000, 0100, 1000\}$$

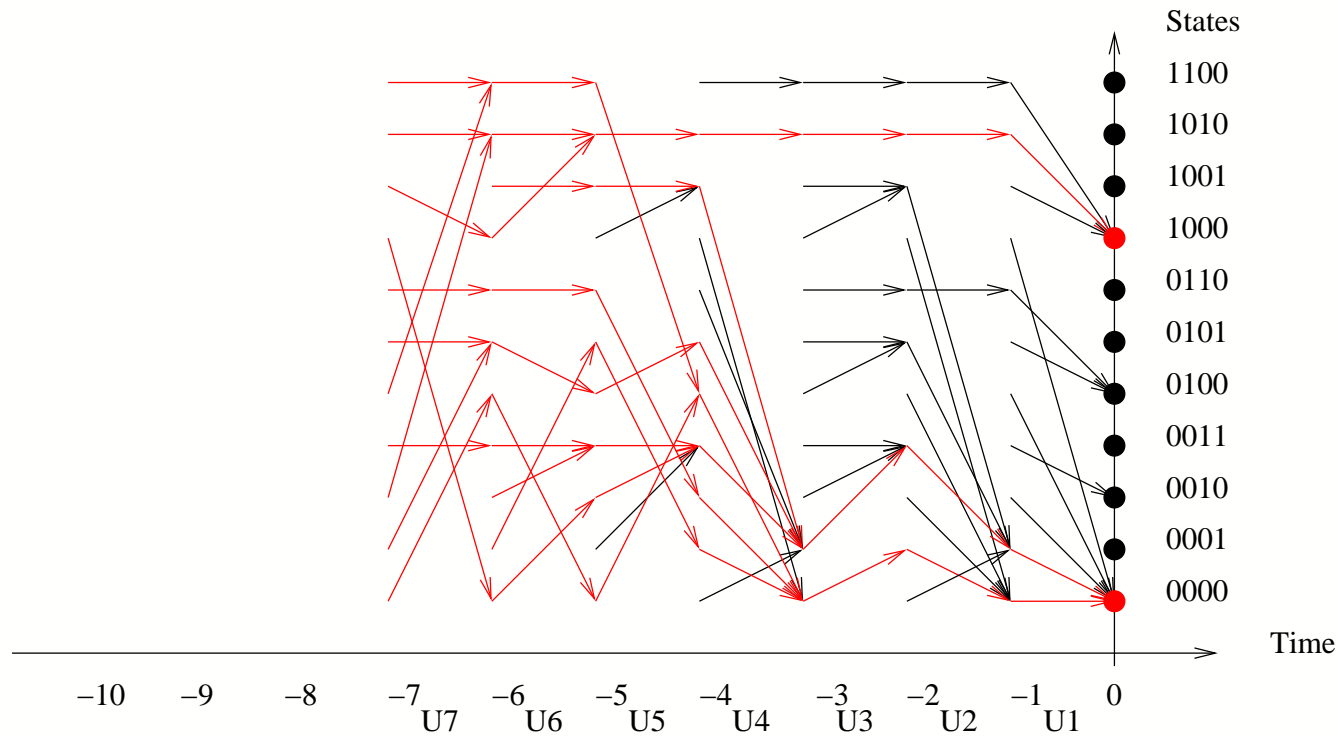
$$\mathcal{Z}_3 = \{0000, 0100, 1000\}$$

$$\mathcal{Z}_4 = \{0000, 1000\}$$

$$\mathcal{Z}_5 = \{0000, 1000\}$$

$$\mathcal{Z}_6 = \{0000, 1000\}$$

# Backward simulation example



$$\mathcal{Z}_0 = \mathcal{X}$$

$$\mathcal{Z}_1 = \{0000, 0010, 0100, 1000\}$$

$$\mathcal{Z}_2 = \{0000, 0100, 1000\}$$

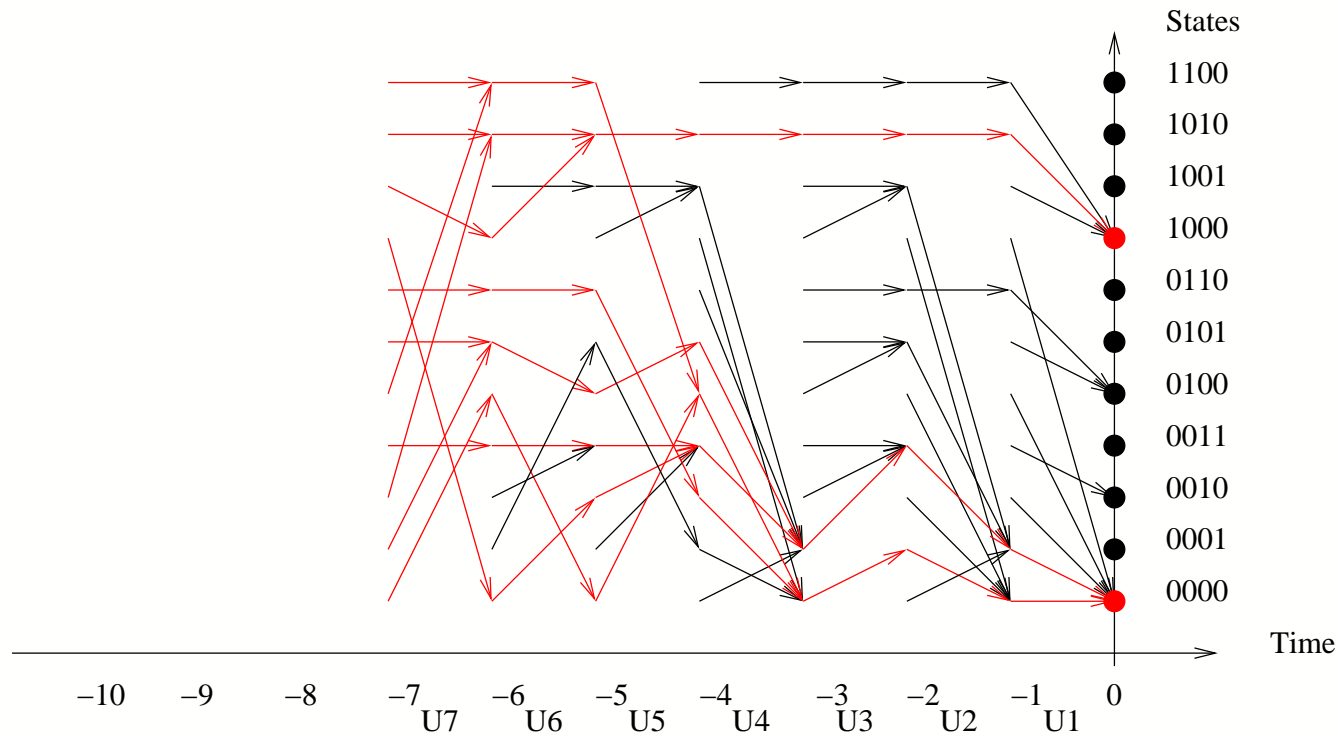
$$\mathcal{Z}_3 = \{0000, 0100, 1000\}$$

$$\mathcal{Z}_4 = \{0000, 1000\}$$

$$\mathcal{Z}_5 = \{0000, 1000\}$$

$$\mathcal{Z}_6 = \{0000, 1000\}$$

# Backward simulation example



$$\mathcal{Z}_0 = \mathcal{X}$$

$$\mathcal{Z}_1 = \{0000, 0010, 0100, 1000\}$$

$$\mathcal{Z}_2 = \{0000, 0100, 1000\}$$

$$\mathcal{Z}_3 = \{0000, 0100, 1000\}$$

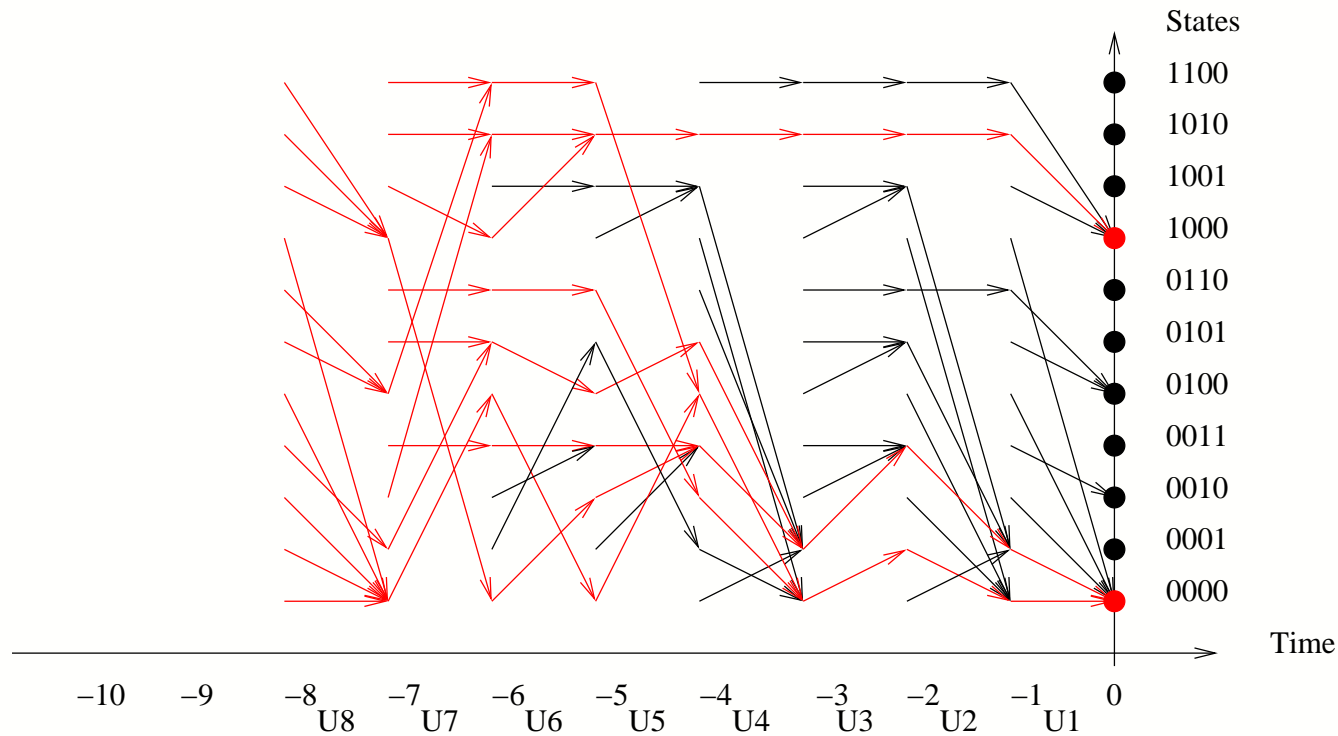
$$\mathcal{Z}_4 = \{0000, 1000\}$$

$$\mathcal{Z}_5 = \{0000, 1000\}$$

$$\mathcal{Z}_6 = \{0000, 1000\}$$

$$\mathcal{Z}_7 = \{0000, 1000\}$$

# Backward simulation example



$$\mathcal{Z}_0 = \mathcal{X}$$

$$\mathcal{Z}_1 = \{0000, 0010, 0100, 1000\}$$

$$\mathcal{Z}_2 = \{0000, 0100, 1000\}$$

$$\mathcal{Z}_3 = \{0000, 0100, 1000\}$$

$$\mathcal{Z}_4 = \{0000, 1000\}$$

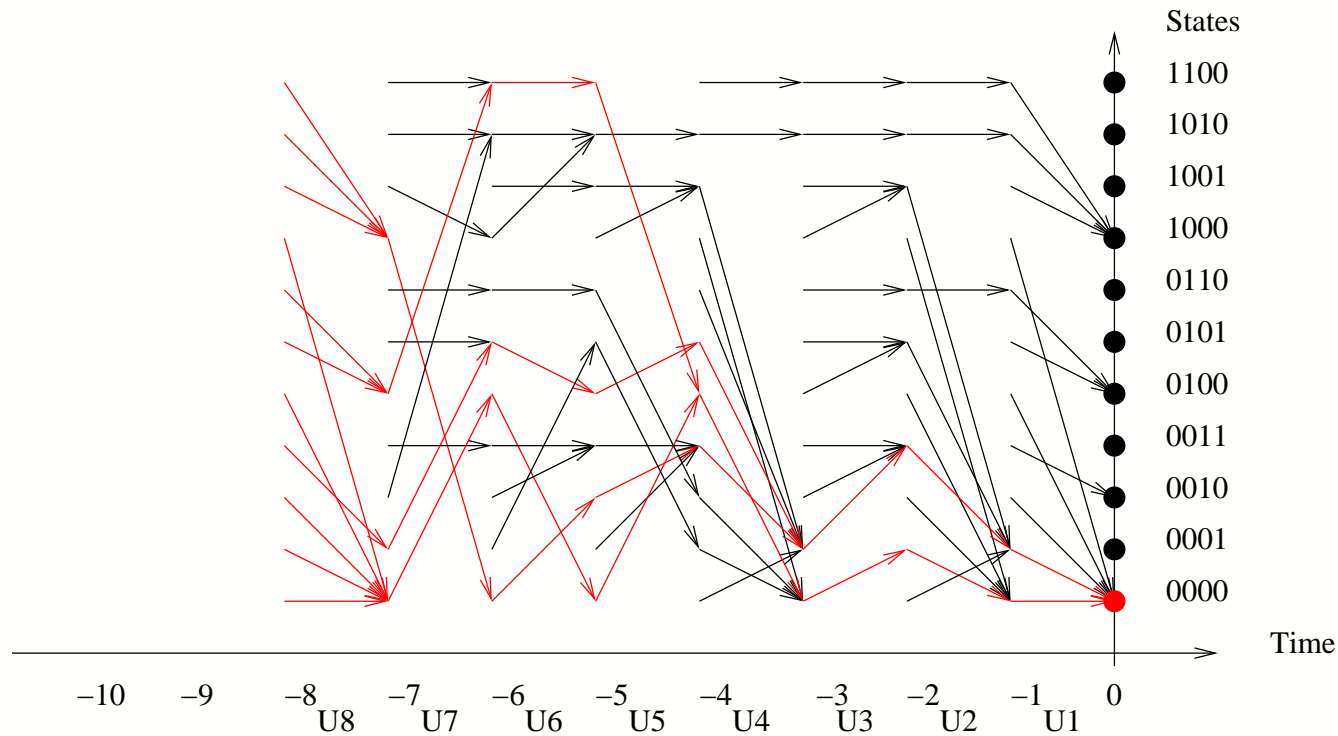
$$\mathcal{Z}_5 = \{0000, 1000\}$$

$$\mathcal{Z}_6 = \{0000, 1000\}$$

$$\mathcal{Z}_7 = \{0000, 1000\}$$



# Backward simulation example



$$\mathcal{Z}_0 = \mathcal{X}$$

$$\mathcal{Z}_1 = \{0000, 0010, 0100, 1000\}$$

$$\mathcal{Z}_2 = \{0000, 0100, 1000\}$$

$$\mathcal{Z}_3 = \{0000, 0100, 1000\}$$

$$\mathcal{Z}_4 = \{0000, 1000\}$$

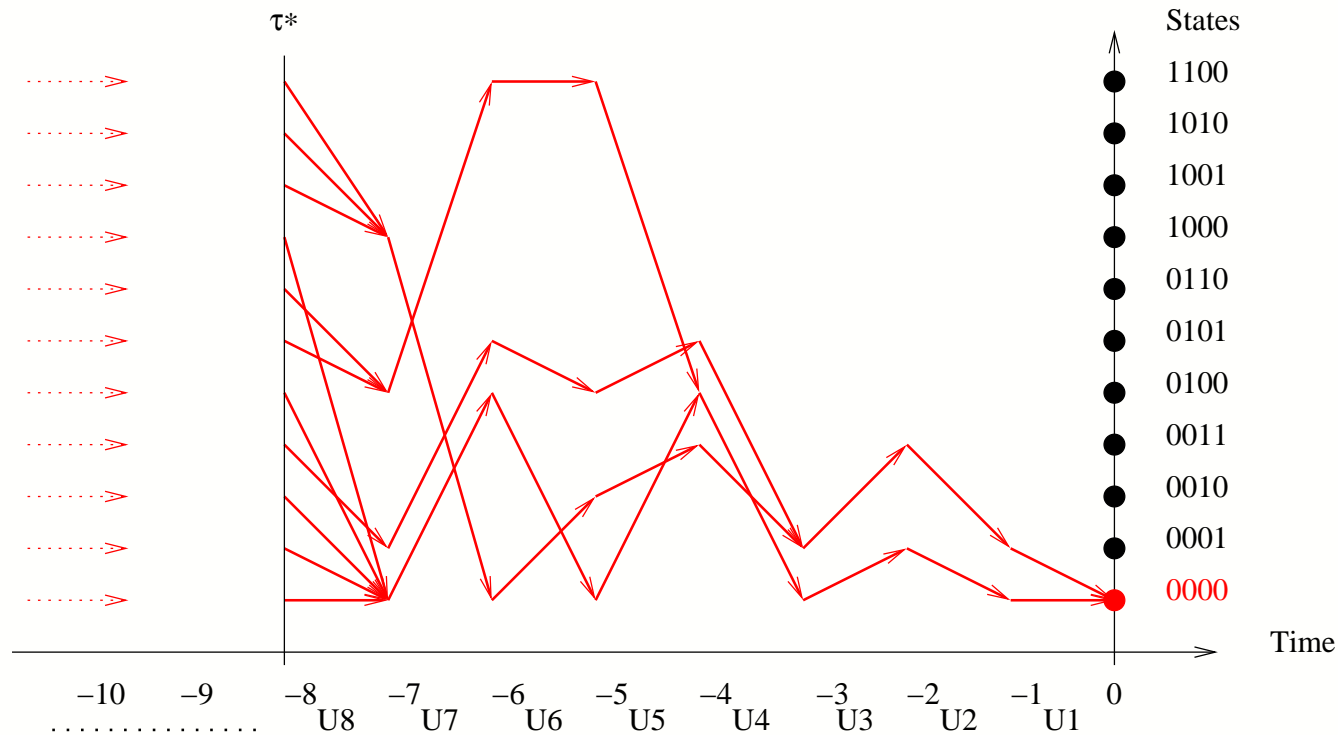
$$\mathcal{Z}_5 = \{0000, 1000\}$$

$$\mathcal{Z}_6 = \{0000, 1000\}$$

$$\mathcal{Z}_7 = \{0000, 1000\}$$

$$\mathcal{Z}_8 = \{0000\}$$

# Backward simulation example



Process stops when  $|\mathcal{Z}_n| = 1$

Stopping time  $\tau^* = 8$

Number of computation of  $\Phi$  (complexity) :  $n \cdot \tau^*$

# Backward coupling simulation

**Proposition 1 (Propp & Wilson)** *If  $\tau^* < +\infty$  a.s. then the returned value is stationary distributed.*

Proof :

$\{Z_n\}_{n \in \mathbb{N}}$  is non-increasing and constant for  $n$  sufficiently large

$$\Phi(\Phi(\cdots(\Phi(\mathcal{X}, U_{-n+1}), \cdots), U_{-1}), U_0) \stackrel{\mathcal{L}}{\sim} \Phi(\Phi(\cdots(\Phi(\mathcal{X}, U_1), \cdots), U_{n-1}), U_n)$$

Remarks :

Stopping times  $\tau$  and  $\tau^*$  have the same law,

$\tau$  depends on  $\Phi$  coding (not only on the transition matrix !)

$\Rightarrow$  optimization problem



# Functional backward coupling

## Idea:

- Directly compute the cost function  $C$
- Same backward scheme
- wait for coupling
- stopping time  $\tau_C^*$

$$\tau_C^* \leq \tau^* \text{ a.s.}$$

⇒ speedup

$$\mathcal{Z}_n^C = C(\Phi(\Phi(\cdots(\Phi(\mathcal{X}, U_{-n+1}), \cdots), U_{-1}), U_0)).$$

potential set of reachable costs at step  $n$

```
for all  $x \in \mathcal{X}$  do  
     $y(x) \leftarrow C(x)$   
end for  
repeat  
     $u \leftarrow \text{Random};$   
    for all  $x \in \mathcal{X}$  do  
         $y(x) \leftarrow y(\Phi(x, u));$   
    end for  
until All  $y(x)$  are equal  
return  $y(x)$ 
```

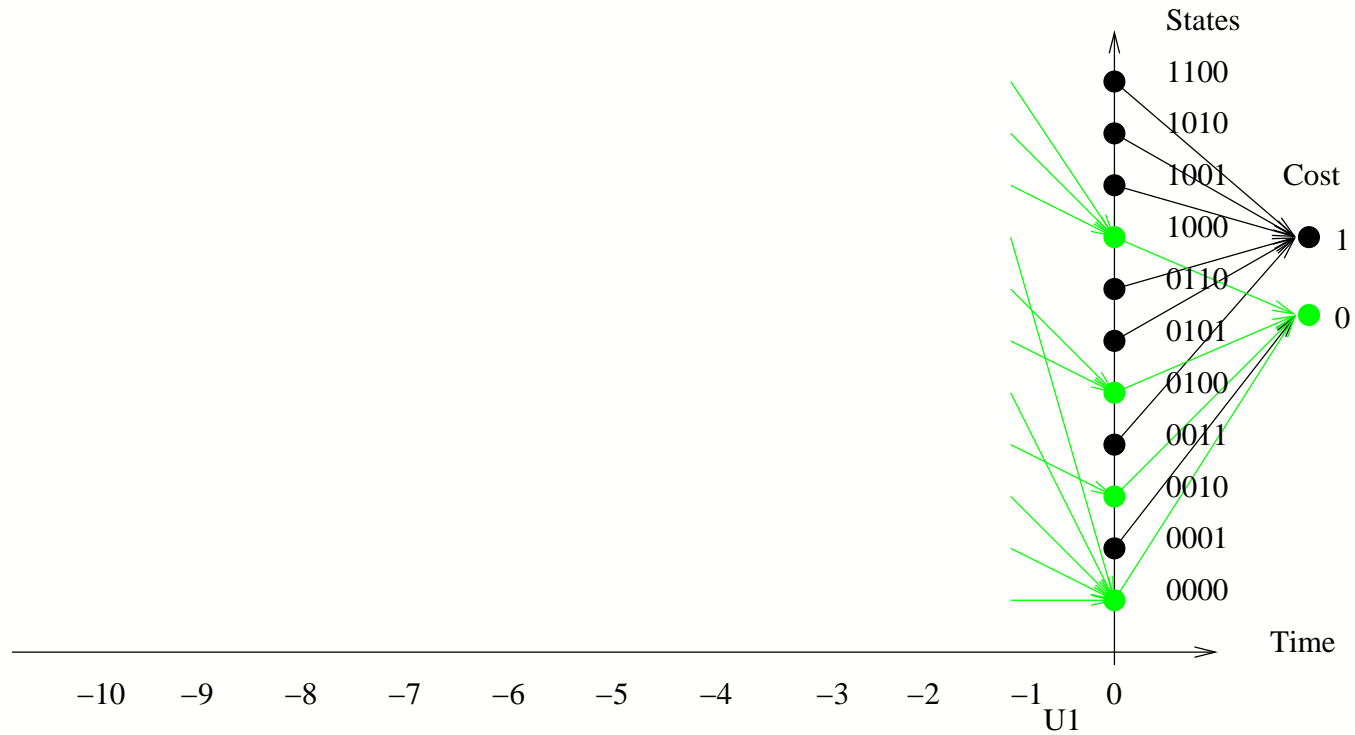


# Backward functional simulation



$$\mathcal{Z}_0^C = \{0, 1\}$$

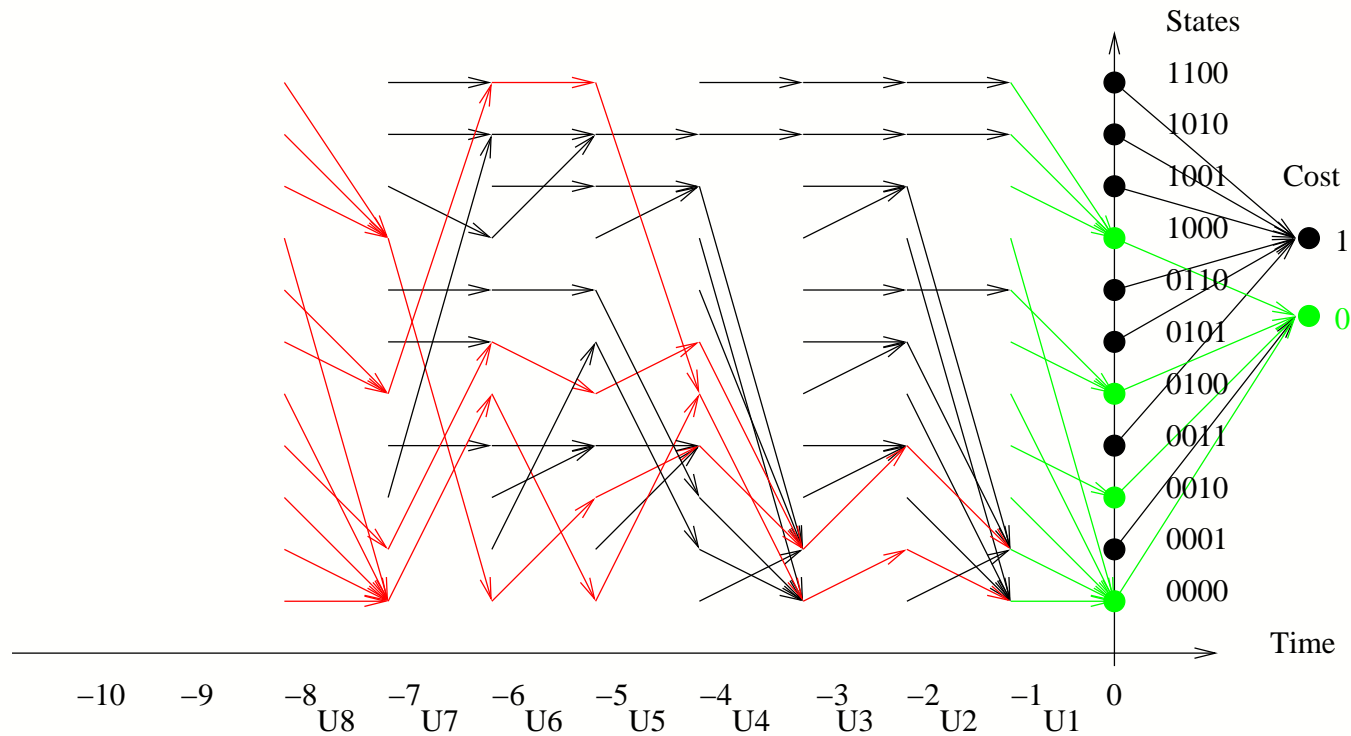
# Backward functional simulation



$$\mathcal{Z}_0^C = \{0, 1\}$$

$$\mathcal{Z}_1^C = \{0\}$$

# Backward functional simulation



$$\mathcal{Z}_0^C = \{0, 1\}$$

$$\mathcal{Z}_1^C = \{0\}$$

$$\tau^* = 8 \quad \tau_C^* = 1$$

**Is the coupling time reduction significant ?**

# Implementation considerations

$\Phi$  coding : choice of a representation

- Sparse description of Markov generators ( $m$  non negative elements)
- Discretization (Uniformization)
- Aliasing table
- Coupling condition





# Examples

- Functionality tests : "random transition matrices"
- Resource sharing model : statistical verification
- Overflow model : sparsity and gain



# Example 1

## Random transition coefficients:

Number of states	10	100	500	1000	3000
Mean coupling time	3.1	4.5	5.3	5.7	6.1
Mean execution time $\mu s$	3	17	170	360	1100

Pentium III 700MHz and 256Mb memory. Sample size 10000.

Remarks:

- **very small coupling time**
- Coefficients : same order of magnitude, aliasing enforces coupling

## Comparison with birth and death process :

Number of states	10	100	500	1000	3000
Mean coupling time	41	557	2850	5680	17000
Mean execution time $\mu s$	28	1800	88177	366000	3.5s

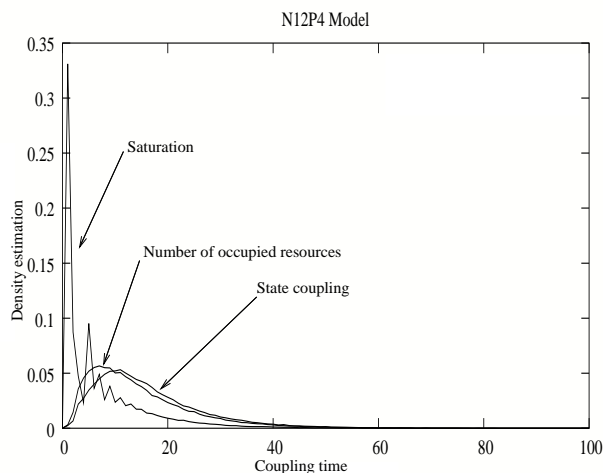
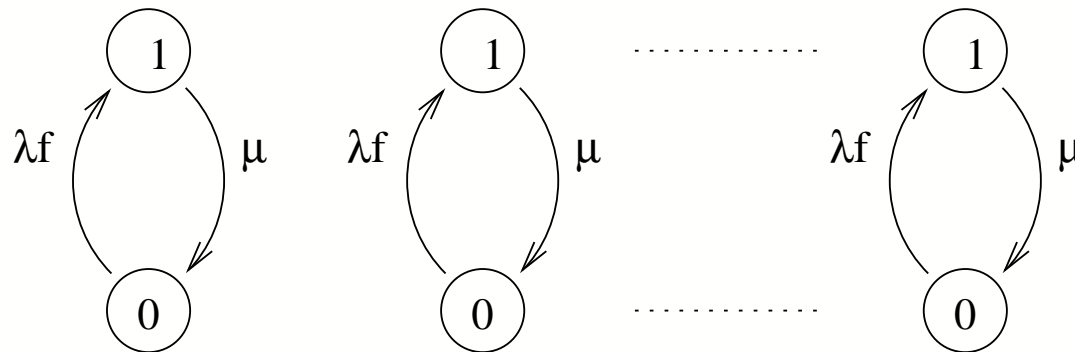
Remarks:

- **large coupling time**
- sparse matrix, large graph diameter



# Resource sharing model

$P$  resources  $N$  users; state  $= (x_1, \dots, x_N)$ ; access constraint  $f = (\sum x_i < P)$   
 product form solution  $\Rightarrow$  statistical validation

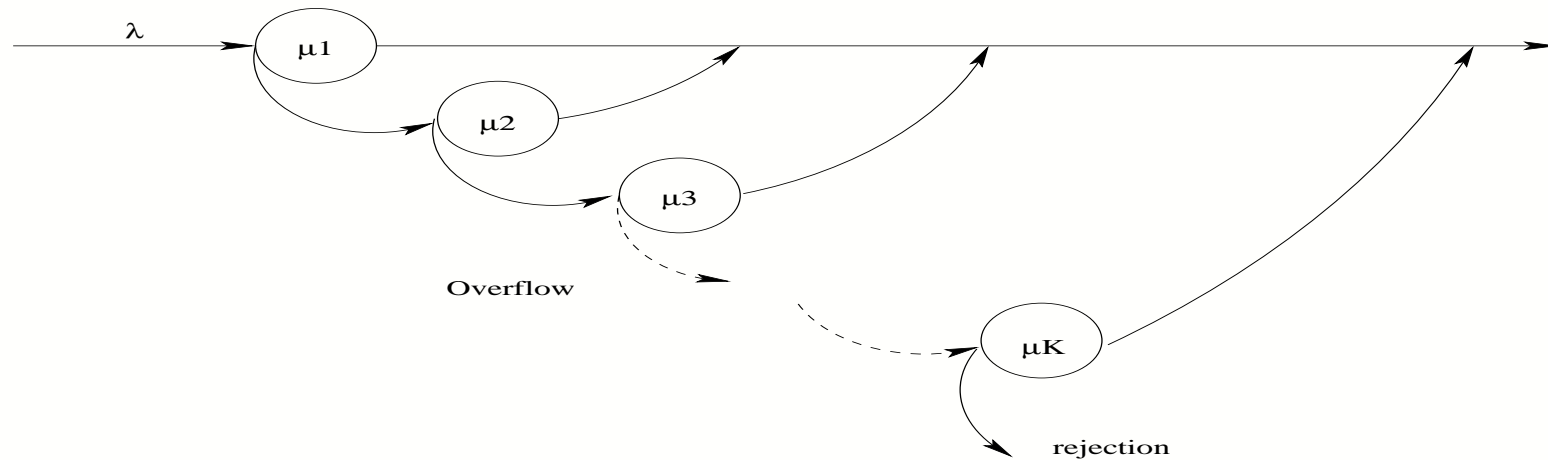


- state coupling
- functional coupling :
- number of occupied resources :
- $\sum_i x_i,$
- saturation :  $(\sum_i x_i = P),$

**High reduction of the functional coupling time**

# Erlang model (1905)

Typical overflow model  $K$  servers



State vector =  $[1, 1, 0, 0, 1, 0, \dots]$

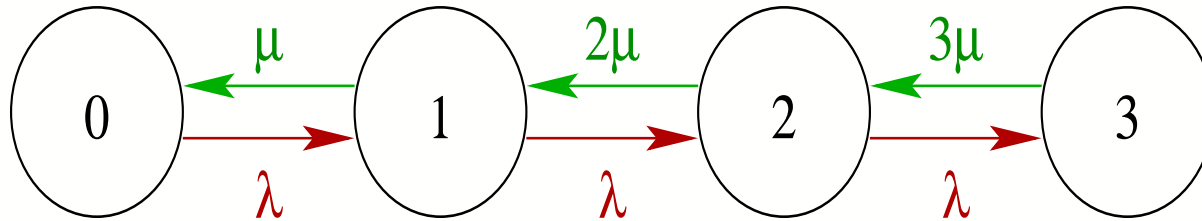
State space size  $2^K$

$\lambda$  input rate

$\mu_i$  service rates

# Erlang model (1905)

$\mu_i = \mu$  aggregation,



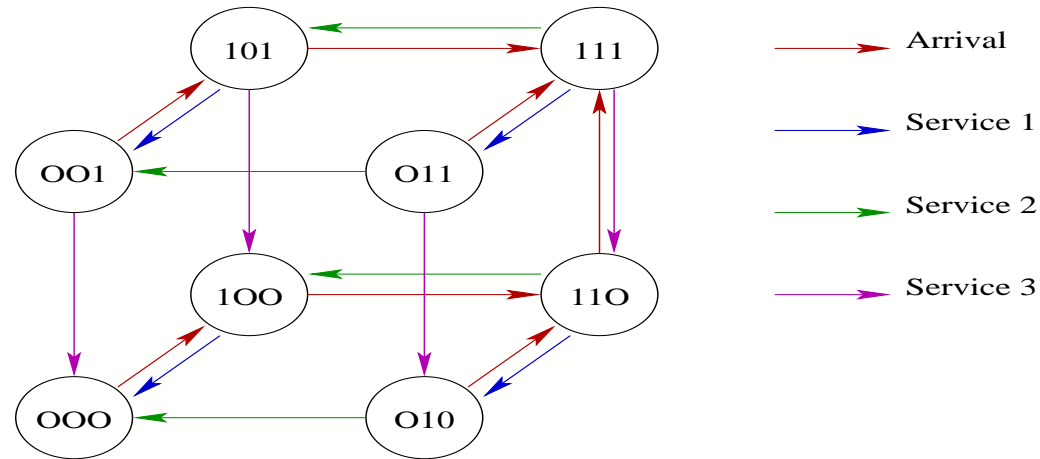
reversibility, product form solution

$$\pi_n = G(K) \frac{1}{n!} \left( \frac{\lambda}{\mu} \right)^n$$

⇒ Erlang formula

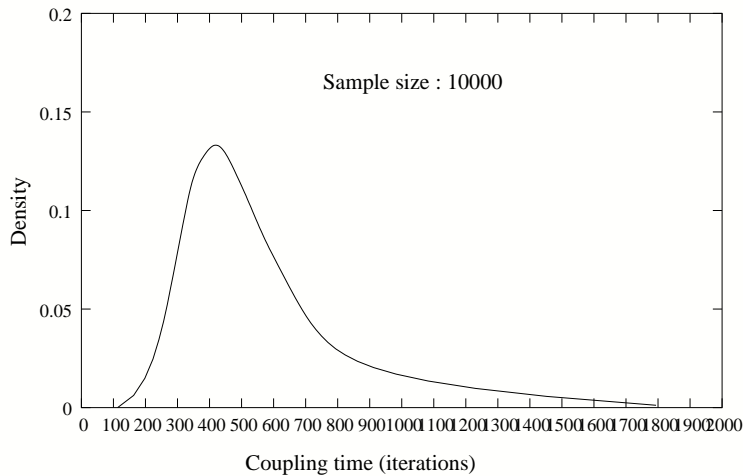
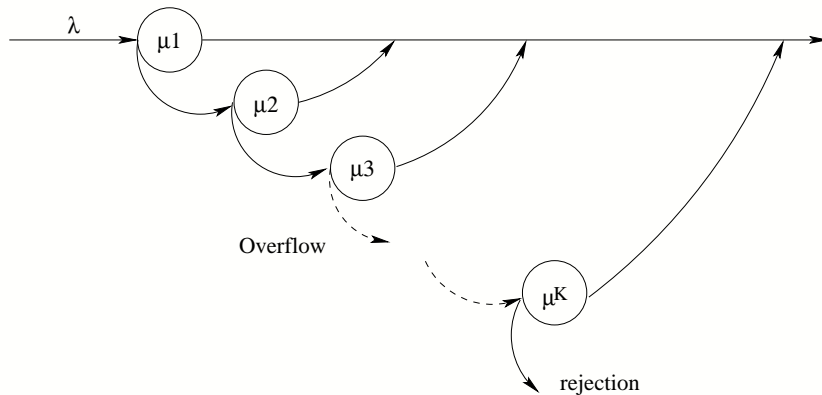
# Erlang model (1905)

$\mu_i$  distincts ???



**non monotonic, non reversible !!!**

# Overflow model

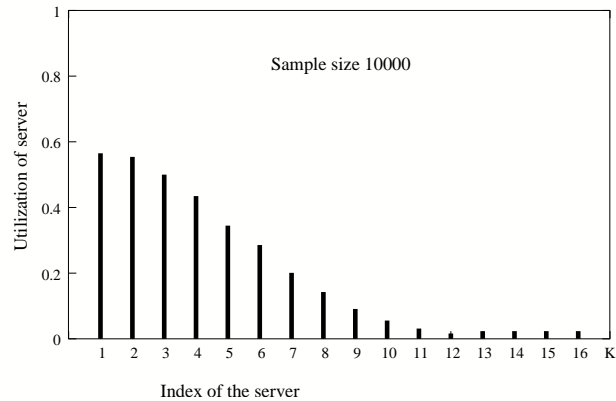


$K$  servers,  
priority on overflows  
input rate  $\lambda$ ,  
different service rate  
state  $(x_1, \dots, x_K)$ ,  $x_i \in \{0, 1\}$ ,  
size  $\sim 130000$   
low diameter  
non product-form structure,

Parameter	Value
minimum	113
maximum	1794
median	465
mean	498
Std	180

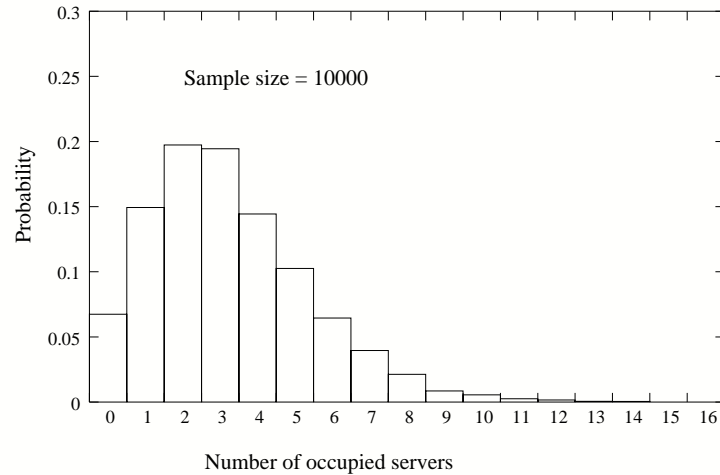
**exponential tail, low mean value**

# Overflow model (2)



Marginal probabilities estimation

$$\mathbb{P}(X_i = 1)$$

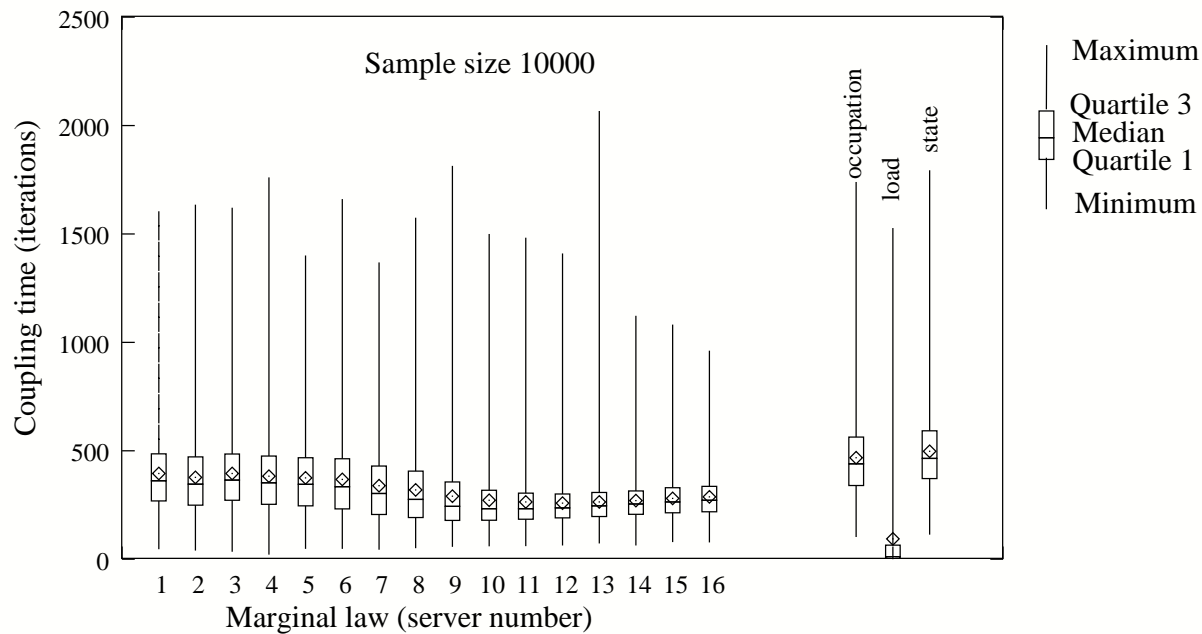


Marginal distribution of the occupied servers





# Overflow model (3)

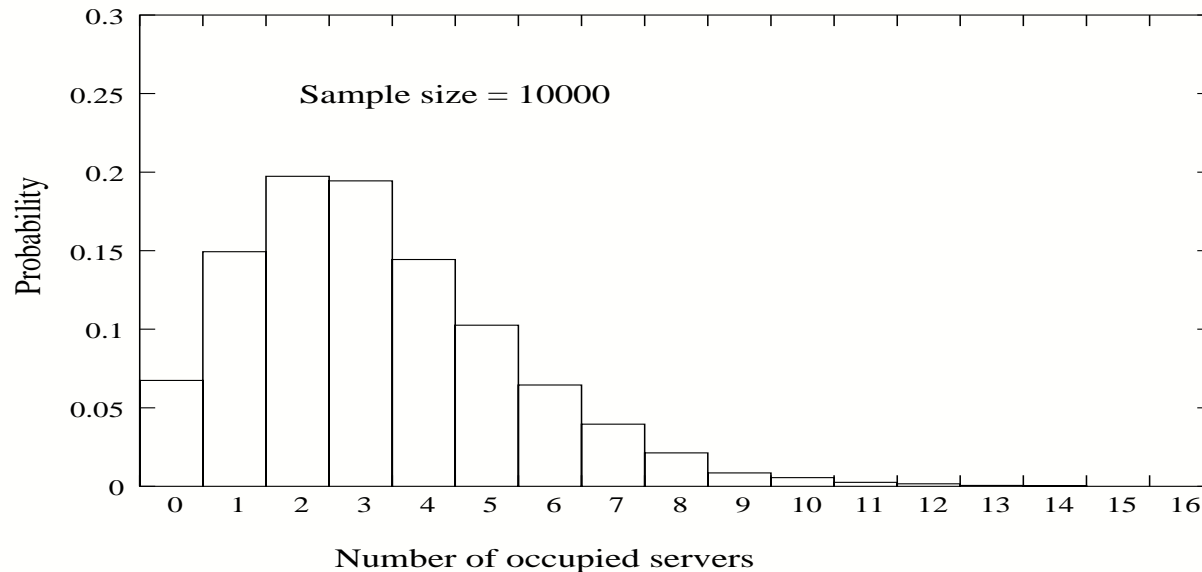


## Functional coupling time

- gain 20% for the first marginals
- utilization : best reduction

# Performances : service number

## Marginal law of occupied servers



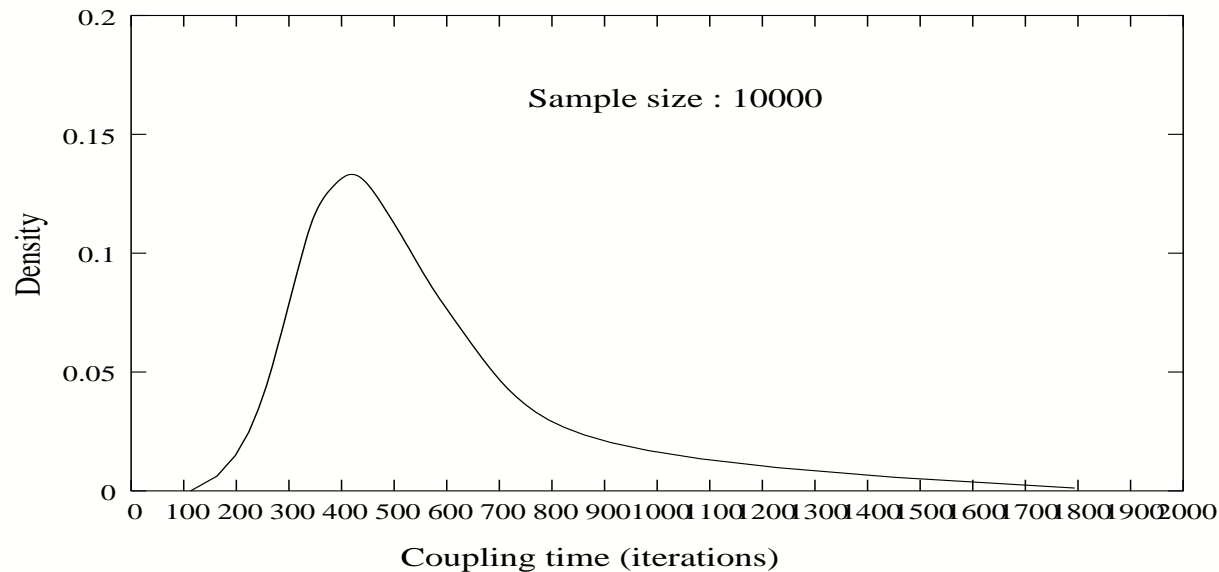
Estimation of the load of the system;

$$\mathbb{P}(N > 13) \leq 10^{-4}$$



# Global coupling

Each iteration operates on 131.072 dimensional vectors

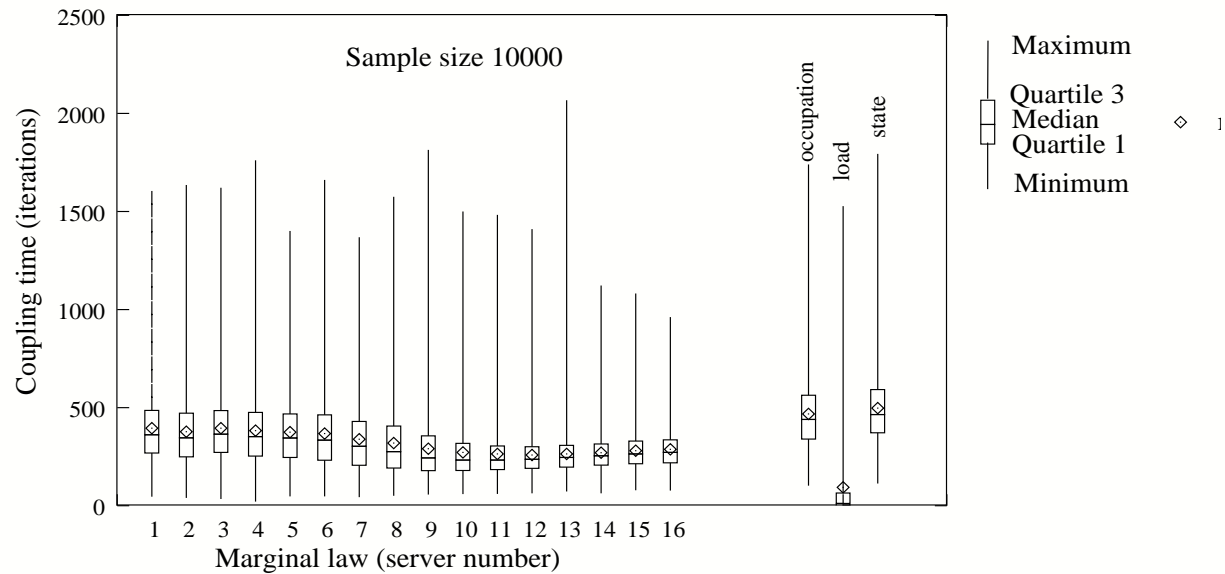


exponential decay



# Coupling time : synthesis

## Functional coupling :



# PSI

Software tool: **PSI : Perfect Simulator**

<http://www-id.imag.fr/Software/PSI/>

Two versions : unix command / with a simple interface

```
psi_alias -i example.marca -o example
```

generates alias tables (example.simu)

```
psi_sample -i example.simu -d sample-size -c example.cost -o  
example
```

example.cost associates to each state its cost

generates samples of costs stationary distributed (example.sample)



# Future works

## Open problems

- relation between  $\phi$  and  $\tau$
- parallel version of the algorithm : block decomposition
- impact of the cost function and aggregation properties

## Technical problems

- interfaces with software modeling tools
- description of the state space (vectors)
- cache optimizations



# Stochastic comparison

A Markov chain is said to be monotonic iff it preserves stochastic ordering on distributions.

Question :

could we find a transition coding function  $\phi$  such that

$$x \leq y \Rightarrow \forall u \quad \phi(x, u) \leq_X \phi(y, u)$$

→ few trajectories are necessary



# From continuous to discrete time

Continuous time Markov chain  
Generator  $Q$  ,  $\Lambda \geq \max_i q_{i,i}$  then

$$P = Id + \frac{1}{\Lambda}Q$$

is the transition matrix of the markov chain of the embedded process driven by a Poisson process with rate  $\Lambda$





# Transition or events

Example Erlang 3 servers : state  $[x_1, x_2, x_3]$ ,  
4 events

departure server 1 :

$$[x_1, x_2, x_3] \rightarrow [(x_1 - 1)^+, x_2, x_3]$$

departure server 2 :

$$[x_1, x_2, x_3] \rightarrow [x_1, (x_2 - 1)^+, x_3]$$

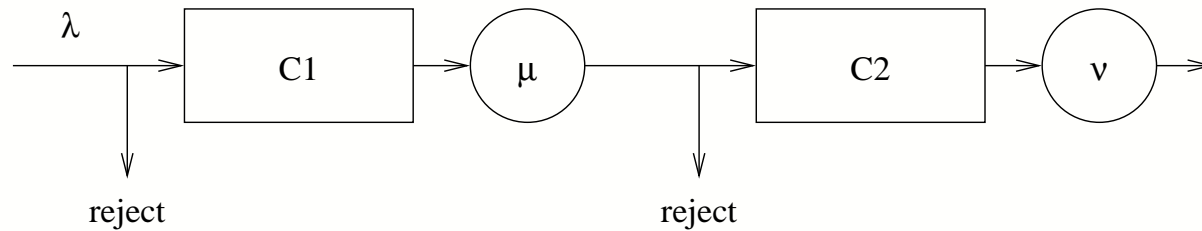
arrival : place 1 in the first non null place

Uniformization by  $\Lambda = \lambda + \mu_1 + \mu_2 + \mu_3$

**$\Rightarrow$  Monotonous process !!!!!**



# Queue or reject



departure server 1 :

$$[x_1, x_2] \rightarrow [(x_1 - 1)^+, \min(x_2 + 1, C2)]$$

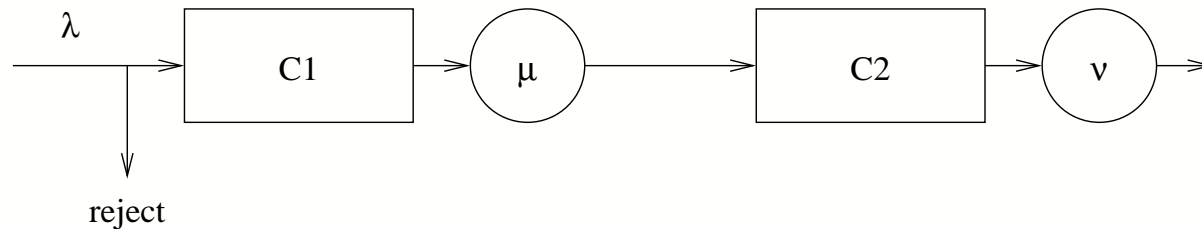
departure server 2 :  $[x_1, x_2] \rightarrow [x_1, (x_2 - 1)^+]$

arrival :  $[x_1, x_2] \rightarrow [\min(x_1 + 1, C1), x_2]$

Uniformization by  $\Lambda = \lambda + \mu + \nu$

**$\Rightarrow$  Monotonous process !!!!!**

# Queue or block



departure server 1 :

$$[x_1, x_2] \rightarrow [x_1 - 1, x_2 + 1] \text{ if } x_1 > 0 \text{ and } x_2 < C2$$

departure server 2 :  $[x_1, x_2] \rightarrow [x_1, (x_2 - 1)^+]$

arrival :  $[x_1, x_2] \rightarrow [\min(x_1 + 1, C1), x_2]$

Uniformization by  $\Lambda = \lambda + \mu + \nu$

**$\Rightarrow$  Monotonous process !!!!!**

# Queueing networks

Event approach : operates on vectors

Monotonicity with rejection, blocking, priorities

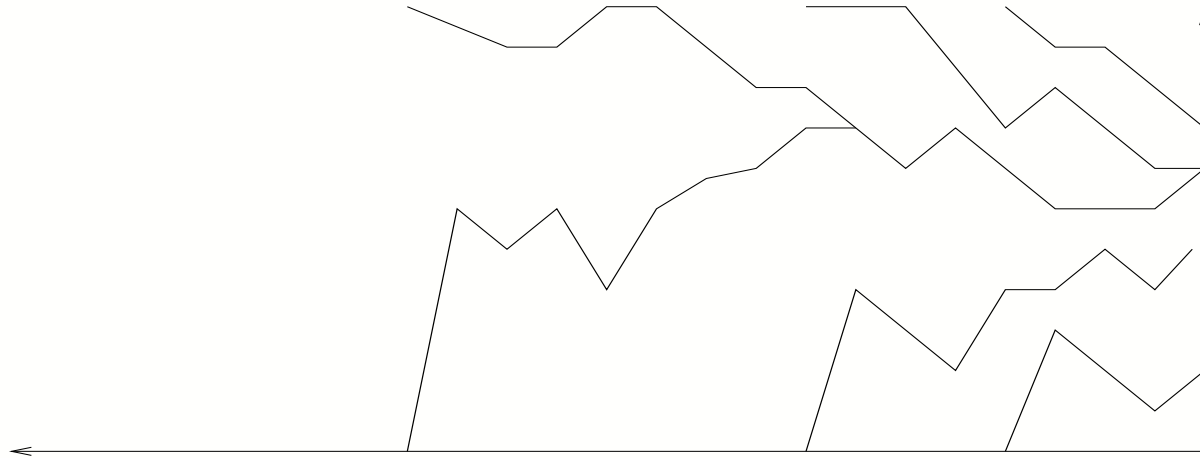
Exercises : find typical non monotonous networks

Generalize to finite number of customers networks (show it)



# Implementation considerations

The global scheme is



needs the storage of events

Propp & Wilson : double period of simulation

Fill : interruptible forward algorithm



# Future works

- Theoretical improvements:
  - deeper understanding of  $\Phi$  properties and the spectrum of the transition matrix
  - evaluation or bounds on the coupling time
- Algorithmic perspectives:
  - building of alias table,
  - transform of alias table,
  - parallelization
- Model based approach :
  - structuration of the matrix : adapted strategies (QN, SAN, GSPN, PA,...)
  - model properties : monotonicity, reversibility,...
- Experimental results
  - find limit models : (ex Birth and death)
  - significant results (depending on the diameter and the coding of the chain)
  - huge models (size  $2^{22}$ )

