

# Fundamental Computer Science

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(inspired by Giorgio Lucarelli)

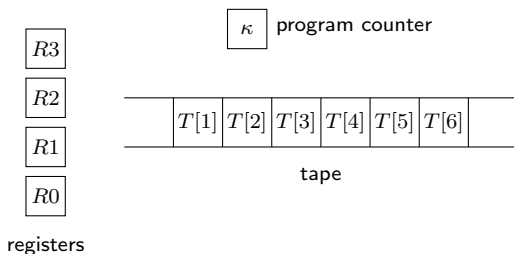
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# Random Access Turing Machines

- ▶ Random Access Memory
  - ▶ access any position of the tape in a single step

# Random Access Turing Machines

- ▶ Random Access Memory
  - ▶ access any position of the tape in a single step
- ▶ we also need:
  - ▶ finite number of *registers* → manipulate addresses of the tape
  - ▶ *program counter* → current **instruction** to execute



- ▶ program: a set of instructions

# Random Access Turing Machines: Instructions set

instruction	operand	semantics
read	$j$	$R_0 \leftarrow T[R_j]$
write	$j$	$T[R_j] \leftarrow R_0$
store	$j$	$R_j \leftarrow R_0$
load	$j$	$R_0 \leftarrow R_j$
load	$= c$	$R_0 = c$
add	$j$	$R_0 \leftarrow R_0 + R_j$
add	$= c$	$R_0 \leftarrow R_0 + c$
sub	$j$	$R_0 \leftarrow \max\{R_0 + R_j, 0\}$
sub	$= c$	$R_0 \leftarrow \max\{R_0 + c, 0\}$
half		$R_0 \leftarrow \lfloor \frac{R_0}{2} \rfloor$
jump	$s$	$\kappa \leftarrow s$
jpos	$s$	if $R_0 > 0$ then $\kappa \leftarrow s$
jzero	$s$	if $R_0 = 0$ then $\kappa \leftarrow s$
halt		$\kappa = 0$

► register  $R_0$ : *accumulator*

# Random Access Turing Machines: Formal definition

A Random Access Turing Machine is a pair  $M = (k, \Pi)$ , where

- ▶  $k > 0$  is the finite number of registers, and
- ▶  $\Pi = (\pi_1, \pi_2, \dots, \pi_p)$  is a finite sequence of instructions (program).

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## Notations

- ▶ the last instruction  $\pi_p$  is always a *halt* instruction
- ▶  $(\kappa; R_0, R_1, \dots, R_{k-1}; T)$ : a **configuration**, where
  - ▶  $\kappa$ : program counter
  - ▶  $R_j, 0 \leq j < k$ : the current value of register  $j$
  - ▶  $T$ : the contents of the tape  
(each  $T[j]$  contains a non-negative integer, i.e.  $T[j] \in \mathbb{N}$ )
- ▶ **halted configuration**:  $\kappa = 0$

# Examples

1: load 1           (1; 0, 5, 3;  $\emptyset$ )  
2: add 2  
3: sub =1  
4: store 1  
5: halt

# Examples

1: load 1             $(1; 0, 5, 3; \emptyset) \vdash (2; 5, 5, 3; \emptyset) \vdash (3; 8, 5, 3; \emptyset) \vdash (4; 7, 5, 3; \emptyset)$   
2: add 2              $\vdash (5; 7, 7, 3; \emptyset) \vdash (0; 7, 7, 3; \emptyset)$   
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# Examples

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5: halt             $R_1 \leftarrow R_2 + R_1 - 1$

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5: halt

$$R_1 \leftarrow R_2 + R_1 - 1$$

1: load 1             $(1; 0, 7; \emptyset)$   
2: jzero 6  
3: sub =3  
4: store 1  
5: jump 2  
6: halt

# Examples

1: load 1             $(1; 0, 5, 3; \emptyset) \vdash (2; 5, 5, 3; \emptyset) \vdash (3; 8, 5, 3; \emptyset) \vdash (4; 7, 5, 3; \emptyset)$   
2: add 2              $\vdash (5; 7, 7, 3; \emptyset) \vdash (0; 7, 7, 3; \emptyset)$   
3: sub =1  
4: store 1  
5: halt

$$R_1 \leftarrow R_2 + R_1 - 1$$

1: load 1             $(1; 0, 7; \emptyset) \vdash (2; 7, 7; \emptyset) \vdash (3; 7, 7; \emptyset) \vdash (4; 4, 7; \emptyset) \vdash (5; 4, 4; \emptyset)$   
2: jzero 6            $\vdash (2; 4, 4; \emptyset) \vdash (3; 4, 4; \emptyset) \vdash (4; 1, 4; \emptyset) \vdash (5; 1, 1; \emptyset)$   
3: sub =3             $\vdash (2; 1, 1; \emptyset) \vdash (3; 1, 1; \emptyset) \vdash (4; 0, 1; \emptyset) \vdash (5; 0, 0; \emptyset)$   
4: store 1  
5: jump 2  
6: halt                $\vdash (2; 0, 0; \emptyset) \vdash (6; 0, 0; \emptyset) \vdash (0; 0, 0; \emptyset)$

# Examples

1: load 1             $(1; 0, 5, 3; \emptyset) \vdash (2; 5, 5, 3; \emptyset) \vdash (3; 8, 5, 3; \emptyset) \vdash (4; 7, 5, 3; \emptyset)$   
2: add 2              $\vdash (5; 7, 7, 3; \emptyset) \vdash (0; 7, 7, 3; \emptyset)$   
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2: jzero 6            $\vdash (2; 4, 4; \emptyset) \vdash (3; 4, 4; \emptyset) \vdash (4; 1, 4; \emptyset) \vdash (5; 1, 1; \emptyset)$   
3: sub =3             $\vdash (2; 1, 1; \emptyset) \vdash (3; 1, 1; \emptyset) \vdash (4; 0, 1; \emptyset) \vdash (5; 0, 0; \emptyset)$   
4: store 1  
5: jump 2  
6: halt

$$\text{while } R_1 > 0 \text{ do } R_1 \leftarrow R_1 - 3$$

## Exercise

- ▶ Write a program for a Random Access Turing Machine that multiplies two integers.

Tip: assume that the initial configuration is  $(1; 0, a_1, a_2, 0; \emptyset)$